

Schemes of Existence of some Types of Functions

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Summary. We prove some useful schemes of existence of real sequences, partial functions from a domain into a domain, partial functions from a set to a set and functions from a domain into a domain. At the beginning we prove some related auxiliary theorems to the article [1].

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The notation and terminology used here are introduced in the following articles: [9], [5], [1], [2], [3], [8], [6], [4], and [7]. We adopt the following convention: x , y will be arbitrary, n , m will denote natural numbers, and r will denote a real number. Next we state four propositions:

- (1) For every n there exists m such that $n = 2 \cdot m$ or $n = 2 \cdot m + 1$.
- (2) For every n there exists m such that $n = 3 \cdot m$ or $n = 3 \cdot m + 1$ or $n = 3 \cdot m + 2$.
- (3) For every n there exists m such that $n = 4 \cdot m$ or $n = 4 \cdot m + 1$ or $n = 4 \cdot m + 2$ or $n = 4 \cdot m + 3$.
- (4) For every n there exists m such that $n = 5 \cdot m$ or $n = 5 \cdot m + 1$ or $n = 5 \cdot m + 2$ or $n = 5 \cdot m + 3$ or $n = 5 \cdot m + 4$.

In this article we present several logical schemes. The scheme *ExRealSubseq* concerns a sequence of real numbers \mathcal{A} , and a unary predicate \mathcal{P} , and states that:

there exists a sequence of real numbers q such that q is a subsequence of \mathcal{A} and for every n holds $\mathcal{P}[q(n)]$ and for every n such that for every r such that $r = \mathcal{A}(n)$ holds $\mathcal{P}[r]$ there exists m such that $\mathcal{A}(n) = q(m)$ provided the following requirement is met:

- for every n there exists m such that $n \leq m$ and $\mathcal{P}[\mathcal{A}(m)]$.

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The scheme *ExRealSeq2* deals with a unary functor \mathcal{F} yielding a real number and a unary functor \mathcal{G} yielding a real number and states that:

there exists a sequence of real numbers s such that for every n holds $s(2 \cdot n) = \mathcal{F}(n)$ and $s(2 \cdot n + 1) = \mathcal{G}(n)$

for all values of the parameters.

The scheme *ExRealSeq3* deals with a unary functor \mathcal{F} yielding a real number, a unary functor \mathcal{G} yielding a real number, and a unary functor \mathcal{H} yielding a real number and states that:

there exists a sequence of real numbers s such that for every n holds $s(3 \cdot n) = \mathcal{F}(n)$ and $s(3 \cdot n + 1) = \mathcal{G}(n)$ and $s(3 \cdot n + 2) = \mathcal{H}(n)$

for all values of the parameters.

The scheme *ExRealSeq4* deals with a unary functor \mathcal{F} yielding a real number, a unary functor \mathcal{G} yielding a real number, a unary functor \mathcal{H} yielding a real number, and a unary functor \mathcal{I} yielding a real number and states that:

there exists a sequence of real numbers s such that for every n holds $s(4 \cdot n) = \mathcal{F}(n)$ and $s(4 \cdot n + 1) = \mathcal{G}(n)$ and $s(4 \cdot n + 2) = \mathcal{H}(n)$ and $s(4 \cdot n + 3) = \mathcal{I}(n)$

for all values of the parameters.

The scheme *ExRealSeq5* deals with a unary functor \mathcal{F} yielding a real number, a unary functor \mathcal{G} yielding a real number, a unary functor \mathcal{H} yielding a real number, a unary functor \mathcal{I} yielding a real number, and a unary functor \mathcal{J} yielding a real number and states that:

there exists a sequence of real numbers s such that for every n holds $s(5 \cdot n) = \mathcal{F}(n)$ and $s(5 \cdot n + 1) = \mathcal{G}(n)$ and $s(5 \cdot n + 2) = \mathcal{H}(n)$ and $s(5 \cdot n + 3) = \mathcal{I}(n)$ and $s(5 \cdot n + 4) = \mathcal{J}(n)$

for all values of the parameters.

The scheme *PartFuncExD2* deals with a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , and two unary predicates \mathcal{P} and \mathcal{Q} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every element c of \mathcal{A} holds $c \in \text{dom } f$ if and only if $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ and for every element c of \mathcal{A} such that $c \in \text{dom } f$ holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$

provided the following condition is met:

- for every element c of \mathcal{A} such that $\mathcal{P}[c]$ holds not $\mathcal{Q}[c]$.

The scheme *PartFuncExD2'* concerns a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , and two unary predicates \mathcal{P} and \mathcal{Q} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every element c of \mathcal{A} holds $c \in \text{dom } f$ if and only if $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ and for every element c of \mathcal{A} such that $c \in \text{dom } f$ holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$

provided the following requirement is met:

- for every element c of \mathcal{A} such that $\mathcal{P}[c]$ and $\mathcal{Q}[c]$ holds $\mathcal{F}(c) = \mathcal{G}(c)$.

The scheme *PartFuncExD2''* deals with a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , and a unary predicate \mathcal{P} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that f is total and for every element c of \mathcal{A} such that $c \in \text{dom } f$ holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if not $\mathcal{P}[c]$, then $f(c) = \mathcal{G}(c)$

for all values of the parameters.

The scheme *PartFuncExD3* deals with a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , a unary functor \mathcal{H} yielding an element of \mathcal{B} , and three unary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every element c of \mathcal{A} holds $c \in \text{dom } f$ if and only if $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ or $\mathcal{R}[c]$ and for every element c of \mathcal{A} such that $c \in \text{dom } f$ holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$ but if $\mathcal{R}[c]$, then $f(c) = \mathcal{H}(c)$

provided the parameters satisfy the following condition:

- for every element c of \mathcal{A} holds if $\mathcal{P}[c]$, then not $\mathcal{Q}[c]$ but if $\mathcal{P}[c]$, then not $\mathcal{R}[c]$ but if $\mathcal{Q}[c]$, then not $\mathcal{R}[c]$.

The scheme *PartFuncExD3'* concerns a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , a unary functor \mathcal{H} yielding an element of \mathcal{B} , and three unary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every element c of \mathcal{A} holds $c \in \text{dom } f$ if and only if $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ or $\mathcal{R}[c]$ and for every element c of \mathcal{A} such that $c \in \text{dom } f$ holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$ but if $\mathcal{R}[c]$, then $f(c) = \mathcal{H}(c)$

provided the following requirement is met:

- for every element c of \mathcal{A} holds if $\mathcal{P}[c]$ and $\mathcal{Q}[c]$, then $\mathcal{F}(c) = \mathcal{G}(c)$ but if $\mathcal{P}[c]$ and $\mathcal{R}[c]$, then $\mathcal{F}(c) = \mathcal{H}(c)$ but if $\mathcal{Q}[c]$ and $\mathcal{R}[c]$, then $\mathcal{G}(c) = \mathcal{H}(c)$.

The scheme *PartFuncExD4* deals with a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , a unary functor \mathcal{H} yielding an element of \mathcal{B} , a unary functor \mathcal{I} yielding an element of \mathcal{B} , and four unary predicates \mathcal{P} , \mathcal{Q} , \mathcal{R} , and \mathcal{S} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every element c of \mathcal{A} holds $c \in \text{dom } f$ if and only if $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ or $\mathcal{R}[c]$ or $\mathcal{S}[c]$ and for every element c of \mathcal{A} such that $c \in \text{dom } f$ holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$ but if $\mathcal{R}[c]$, then $f(c) = \mathcal{H}(c)$ but if $\mathcal{S}[c]$, then $f(c) = \mathcal{I}(c)$

provided the parameters satisfy the following condition:

- for every element c of \mathcal{A} holds if $\mathcal{P}[c]$, then not $\mathcal{Q}[c]$ but if $\mathcal{P}[c]$, then not $\mathcal{R}[c]$ but if $\mathcal{P}[c]$, then not $\mathcal{S}[c]$ but if $\mathcal{Q}[c]$, then not $\mathcal{R}[c]$ but if $\mathcal{Q}[c]$, then not $\mathcal{S}[c]$ but if $\mathcal{R}[c]$, then not $\mathcal{S}[c]$.

The scheme *PartFuncExS2* deals with a set \mathcal{A} , a set \mathcal{B} , a unary functor \mathcal{F} , a unary functor \mathcal{G} , and two unary predicates \mathcal{P} and \mathcal{Q} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every x holds $x \in \text{dom } f$ if and only if $x \in \mathcal{A}$ but $\mathcal{P}[x]$ or $\mathcal{Q}[x]$ and for every x such that $x \in \text{dom } f$ holds if $\mathcal{P}[x]$, then $f(x) = \mathcal{F}(x)$ but if $\mathcal{Q}[x]$, then $f(x) = \mathcal{G}(x)$

provided the parameters satisfy the following conditions:

- for every x such that $x \in \mathcal{A}$ holds if $\mathcal{P}[x]$, then not $\mathcal{Q}[x]$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{P}[x]$ holds $\mathcal{F}(x) \in \mathcal{B}$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{Q}[x]$ holds $\mathcal{G}(x) \in \mathcal{B}$.

The scheme *PartFuncExS3* deals with a set \mathcal{A} , a set \mathcal{B} , a unary functor \mathcal{F} , a unary functor \mathcal{G} , a unary functor \mathcal{H} , and three unary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every x holds $x \in \text{dom } f$ if and only if $x \in \mathcal{A}$ but $\mathcal{P}[x]$ or $\mathcal{Q}[x]$ or $\mathcal{R}[x]$ and for every x such that $x \in \text{dom } f$ holds if $\mathcal{P}[x]$, then $f(x) = \mathcal{F}(x)$ but if $\mathcal{Q}[x]$, then $f(x) = \mathcal{G}(x)$ but if $\mathcal{R}[x]$, then $f(x) = \mathcal{H}(x)$

provided the parameters meet the following conditions:

- for every x such that $x \in \mathcal{A}$ holds if $\mathcal{P}[x]$, then not $\mathcal{Q}[x]$ but if $\mathcal{P}[x]$, then not $\mathcal{R}[x]$ but if $\mathcal{Q}[x]$, then not $\mathcal{R}[x]$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{P}[x]$ holds $\mathcal{F}(x) \in \mathcal{B}$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{Q}[x]$ holds $\mathcal{G}(x) \in \mathcal{B}$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{R}[x]$ holds $\mathcal{H}(x) \in \mathcal{B}$.

The scheme *PartFuncExS4* deals with a set \mathcal{A} , a set \mathcal{B} , a unary functor \mathcal{F} , a unary functor \mathcal{G} , a unary functor \mathcal{H} , a unary functor \mathcal{I} , and four unary predicates \mathcal{P} , \mathcal{Q} , \mathcal{R} , and \mathcal{S} , and states that:

there exists a partial function f from \mathcal{A} to \mathcal{B} such that for every x holds $x \in \text{dom } f$ if and only if $x \in \mathcal{A}$ but $\mathcal{P}[x]$ or $\mathcal{Q}[x]$ or $\mathcal{R}[x]$ or $\mathcal{S}[x]$ and for every x such that $x \in \text{dom } f$ holds if $\mathcal{P}[x]$, then $f(x) = \mathcal{F}(x)$ but if $\mathcal{Q}[x]$, then $f(x) = \mathcal{G}(x)$ but if $\mathcal{R}[x]$, then $f(x) = \mathcal{H}(x)$ but if $\mathcal{S}[x]$, then $f(x) = \mathcal{I}(x)$

provided the parameters meet the following requirements:

- for every x such that $x \in \mathcal{A}$ holds if $\mathcal{P}[x]$, then not $\mathcal{Q}[x]$ but if $\mathcal{P}[x]$, then not $\mathcal{R}[x]$ but if $\mathcal{P}[x]$, then not $\mathcal{S}[x]$ but if $\mathcal{Q}[x]$, then not $\mathcal{R}[x]$ but if $\mathcal{Q}[x]$, then not $\mathcal{S}[x]$ but if $\mathcal{R}[x]$, then not $\mathcal{S}[x]$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{P}[x]$ holds $\mathcal{F}(x) \in \mathcal{B}$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{Q}[x]$ holds $\mathcal{G}(x) \in \mathcal{B}$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{R}[x]$ holds $\mathcal{H}(x) \in \mathcal{B}$,
- for every x such that $x \in \mathcal{A}$ and $\mathcal{S}[x]$ holds $\mathcal{I}(x) \in \mathcal{B}$.

The scheme *PartFuncExC-D2* concerns a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a non-empty set \mathcal{C} , a binary functor \mathcal{F} yielding an element of \mathcal{C} , a binary functor \mathcal{G} yielding an element of \mathcal{C} , and two binary predicates \mathcal{P} and \mathcal{Q} , and states that:

there exists a partial function f from $\{ \mathcal{A}, \mathcal{B} \}$ to \mathcal{C} such that for every element c of \mathcal{A} and for every element d of \mathcal{B} holds $\langle c, d \rangle \in \text{dom } f$ if and only if $\mathcal{P}[c, d]$ or $\mathcal{Q}[c, d]$ and for every element c of \mathcal{A} and for every element d of \mathcal{B} such that $\langle c, d \rangle \in \text{dom } f$ holds if $\mathcal{P}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{F}(c, d)$ but if $\mathcal{Q}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{G}(c, d)$

provided the parameters meet the following requirement:

- for every element c of \mathcal{A} and for every element d of \mathcal{B} such that $\mathcal{P}[c, d]$ holds not $\mathcal{Q}[c, d]$.

The scheme *PartFuncExC_D3* concerns a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a non-empty set \mathcal{C} , a binary functor \mathcal{F} yielding an element of \mathcal{C} , a binary functor \mathcal{G} yielding an element of \mathcal{C} , a binary functor \mathcal{H} yielding an element of \mathcal{C} , and three binary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a partial function f from $[\mathcal{A}, \mathcal{B}]$ to \mathcal{C} such that for every element c of \mathcal{A} and for every element d of \mathcal{B} holds $\langle c, d \rangle \in \text{dom } f$ if and only if $\mathcal{P}[c, d]$ or $\mathcal{Q}[c, d]$ or $\mathcal{R}[c, d]$ and for every element c of \mathcal{A} and for every element r of \mathcal{B} such that $\langle c, r \rangle \in \text{dom } f$ holds if $\mathcal{P}[c, r]$, then $f(\langle c, r \rangle) = \mathcal{F}(c, r)$ but if $\mathcal{Q}[c, r]$, then $f(\langle c, r \rangle) = \mathcal{G}(c, r)$ but if $\mathcal{R}[c, r]$, then $f(\langle c, r \rangle) = \mathcal{H}(c, r)$

provided the following requirement is met:

- for every element c of \mathcal{A} and for every element s of \mathcal{B} holds if $\mathcal{P}[c, s]$, then not $\mathcal{Q}[c, s]$ but if $\mathcal{P}[c, s]$, then not $\mathcal{R}[c, s]$ but if $\mathcal{Q}[c, s]$, then not $\mathcal{R}[c, s]$.

The scheme *PartFuncExC_S2* concerns a set \mathcal{A} , a set \mathcal{B} , a set \mathcal{C} , a binary functor \mathcal{F} , a binary functor \mathcal{G} , and two binary predicates \mathcal{P} and \mathcal{Q} , and states that:

there exists a partial function f from $[\mathcal{A}, \mathcal{B}]$ to \mathcal{C} such that for all x, y holds $\langle x, y \rangle \in \text{dom } f$ if and only if $x \in \mathcal{A}$ and $y \in \mathcal{B}$ but $\mathcal{P}[x, y]$ or $\mathcal{Q}[x, y]$ and for all x, y such that $\langle x, y \rangle \in \text{dom } f$ holds if $\mathcal{P}[x, y]$, then $f(\langle x, y \rangle) = \mathcal{F}(x, y)$ but if $\mathcal{Q}[x, y]$, then $f(\langle x, y \rangle) = \mathcal{G}(x, y)$

provided the following conditions are met:

- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ holds if $\mathcal{P}[x, y]$, then not $\mathcal{Q}[x, y]$,
- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ and $\mathcal{P}[x, y]$ holds $\mathcal{F}(x, y) \in \mathcal{C}$,
- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ and $\mathcal{Q}[x, y]$ holds $\mathcal{G}(x, y) \in \mathcal{C}$.

The scheme *PartFuncExC_S3* concerns a set \mathcal{A} , a set \mathcal{B} , a set \mathcal{C} , a binary functor \mathcal{F} , a binary functor \mathcal{G} , a binary functor \mathcal{H} , and three binary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a partial function f from $[\mathcal{A}, \mathcal{B}]$ to \mathcal{C} such that for all x, y holds $\langle x, y \rangle \in \text{dom } f$ if and only if $x \in \mathcal{A}$ and $y \in \mathcal{B}$ but $\mathcal{P}[x, y]$ or $\mathcal{Q}[x, y]$ or $\mathcal{R}[x, y]$ and for all x, y such that $\langle x, y \rangle \in \text{dom } f$ holds if $\mathcal{P}[x, y]$, then $f(\langle x, y \rangle) = \mathcal{F}(x, y)$ but if $\mathcal{Q}[x, y]$, then $f(\langle x, y \rangle) = \mathcal{G}(x, y)$ but if $\mathcal{R}[x, y]$, then $f(\langle x, y \rangle) = \mathcal{H}(x, y)$

provided the following conditions are met:

- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ holds if $\mathcal{P}[x, y]$, then not $\mathcal{Q}[x, y]$ but if $\mathcal{P}[x, y]$, then not $\mathcal{R}[x, y]$ but if $\mathcal{Q}[x, y]$, then not $\mathcal{R}[x, y]$,
- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ holds if $\mathcal{P}[x, y]$, then $\mathcal{F}(x, y) \in \mathcal{C}$,
- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ holds if $\mathcal{Q}[x, y]$, then $\mathcal{G}(x, y) \in \mathcal{C}$,
- for all x, y such that $x \in \mathcal{A}$ and $y \in \mathcal{B}$ holds if $\mathcal{R}[x, y]$, then $\mathcal{H}(x, y) \in \mathcal{C}$.

The scheme *ExFuncD3* concerns a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element

of \mathcal{B} , a unary functor \mathcal{H} yielding an element of \mathcal{B} , and three unary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a function f from \mathcal{A} into \mathcal{B} such that for every element c of \mathcal{A} holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$ but if $\mathcal{R}[c]$, then $f(c) = \mathcal{H}(c)$

provided the parameters satisfy the following conditions:

- for every element c of \mathcal{A} holds if $\mathcal{P}[c]$, then not $\mathcal{Q}[c]$ but if $\mathcal{P}[c]$, then not $\mathcal{R}[c]$ but if $\mathcal{Q}[c]$, then not $\mathcal{R}[c]$,
- for every element c of \mathcal{A} holds $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ or $\mathcal{R}[c]$.

The scheme *ExFuncD4* concerns a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a unary functor \mathcal{F} yielding an element of \mathcal{B} , a unary functor \mathcal{G} yielding an element of \mathcal{B} , a unary functor \mathcal{H} yielding an element of \mathcal{B} , a unary functor \mathcal{I} yielding an element of \mathcal{B} , and four unary predicates \mathcal{P} , \mathcal{Q} , \mathcal{R} , and \mathcal{S} , and states that:

there exists a function f from \mathcal{A} into \mathcal{B} such that for every element c of \mathcal{A} holds if $\mathcal{P}[c]$, then $f(c) = \mathcal{F}(c)$ but if $\mathcal{Q}[c]$, then $f(c) = \mathcal{G}(c)$ but if $\mathcal{R}[c]$, then $f(c) = \mathcal{H}(c)$ but if $\mathcal{S}[c]$, then $f(c) = \mathcal{I}(c)$

provided the following conditions are met:

- for every element c of \mathcal{A} holds if $\mathcal{P}[c]$, then not $\mathcal{Q}[c]$ but if $\mathcal{P}[c]$, then not $\mathcal{R}[c]$ but if $\mathcal{P}[c]$, then not $\mathcal{S}[c]$ but if $\mathcal{Q}[c]$, then not $\mathcal{R}[c]$ but if $\mathcal{Q}[c]$, then not $\mathcal{S}[c]$ but if $\mathcal{R}[c]$, then not $\mathcal{S}[c]$,
- for every element c of \mathcal{A} holds $\mathcal{P}[c]$ or $\mathcal{Q}[c]$ or $\mathcal{R}[c]$ or $\mathcal{S}[c]$.

The scheme *FuncExC.D2* deals with a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a non-empty set \mathcal{C} , a binary functor \mathcal{F} yielding an element of \mathcal{C} , a binary functor \mathcal{G} yielding an element of \mathcal{C} , and a binary predicate \mathcal{P} , and states that:

there exists a function f from $[\mathcal{A}, \mathcal{B}]$ into \mathcal{C} such that for every element c of \mathcal{A} and for every element d of \mathcal{B} such that $\langle c, d \rangle \in \text{dom } f$ holds if $\mathcal{P}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{F}(c, d)$ but if not $\mathcal{P}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{G}(c, d)$

for all values of the parameters.

The scheme *FuncExC.D3* deals with a non-empty set \mathcal{A} , a non-empty set \mathcal{B} , a non-empty set \mathcal{C} , a binary functor \mathcal{F} yielding an element of \mathcal{C} , a binary functor \mathcal{G} yielding an element of \mathcal{C} , a binary functor \mathcal{H} yielding an element of \mathcal{C} , and three binary predicates \mathcal{P} , \mathcal{Q} , and \mathcal{R} , and states that:

there exists a function f from $[\mathcal{A}, \mathcal{B}]$ into \mathcal{C} such that for every element c of \mathcal{A} and for every element d of \mathcal{B} holds $\langle c, d \rangle \in \text{dom } f$ if and only if $\mathcal{P}[c, d]$ or $\mathcal{Q}[c, d]$ or $\mathcal{R}[c, d]$ and for every element c of \mathcal{A} and for every element d of \mathcal{B} such that $\langle c, d \rangle \in \text{dom } f$ holds if $\mathcal{P}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{F}(c, d)$ but if $\mathcal{Q}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{G}(c, d)$ but if $\mathcal{R}[c, d]$, then $f(\langle c, d \rangle) = \mathcal{H}(c, d)$

provided the parameters have the following properties:

- for every element c of \mathcal{A} and for every element d of \mathcal{B} holds if $\mathcal{P}[c, d]$, then not $\mathcal{Q}[c, d]$ but if $\mathcal{P}[c, d]$, then not $\mathcal{R}[c, d]$ but if $\mathcal{Q}[c, d]$, then not $\mathcal{R}[c, d]$,
- for every element c of \mathcal{A} and for every element d of \mathcal{B} holds $\mathcal{P}[c, d]$ or $\mathcal{Q}[c, d]$ or $\mathcal{R}[c, d]$.

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