

Vertex Sequences Induced by Chains ¹

Yatsuka Nakamura
Shinshu University
Department of Information Engineering
Nagano, Japan

Piotr Rudnicki
University of Alberta
Department of Computing Science
Edmonton, Alberta, Canada

Summary. In the three preliminary sections to the article we define two operations on finite sequences which seem to be of general interest. The first is the *cut* operation that extracts a contiguous chunk of a finite sequence from a position to a position. The second operation is a glueing catenation that given two finite sequences catenates them with removal of the first element of the second sequence. The main topic of the article is to define an operation which for a given chain in a graph returns the sequence of vertices through which the chain passes. We define the exact conditions when such an operation is uniquely definable. This is done with the help of the so called two-valued alternating finite sequences. We also prove theorems about the existence of simple chains which are subchains of a given chain. In order to do this we define the notion of a finite subsequence of a typed finite sequence.

MML Identifier: GRAPH_2.

The articles [16], [20], [9], [21], [6], [7], [4], [5], [19], [15], [8], [3], [1], [14], [10], [11], [2], [18], [17], [12], and [13] provide the notation and terminology for this paper.

¹This work was partially supported by Shinshu Endowment for Information Science, NSERC Grant OGP9207 and JSTF award 651-93-S009.

1. PRELIMINARIES

We adopt the following convention: p, q are finite sequences, X, Y are sets, and i, k, l, m, n, r are natural numbers.

The scheme *FinSegRng* deals with natural numbers \mathcal{A}, \mathcal{B} , a unary functor \mathcal{F} yielding a set, and a unary predicate \mathcal{P} , and states that:

$\{\mathcal{F}(i) : \mathcal{A} \leq i \wedge i \leq \mathcal{B} \wedge \mathcal{P}[i]\}$ is finite

for all values of the parameters.

One can prove the following propositions:

- (1) $m + 1 \leq k$ and $k \leq n$ iff there exists a natural number i such that $m \leq i$ and $i < n$ and $k = i + 1$.
- (2) If $q = p \upharpoonright \text{Seg } n$, then $\text{len } q \leq \text{len } p$ and for every i such that $1 \leq i$ and $i \leq \text{len } q$ holds $p(i) = q(i)$.
- (3) If $X \subseteq \text{Seg } k$ and $Y \subseteq \text{dom Sgm } X$, then $\text{Sgm } X \cdot \text{Sgm } Y = \text{Sgm rng}(\text{Sgm } X \upharpoonright Y)$.
- (4) For all natural numbers m, n holds $\overline{\{k : m \leq k \wedge k \leq m + n\}} = n + 1$.
- (5) For every l such that $1 \leq l$ and $l \leq n$ holds $(\text{Sgm}\{k_1 : k_1 \text{ ranges over natural numbers, } m + 1 \leq k_1 \wedge k_1 \leq m + n\})(l) = m + l$.

2. THE CUT OPERATION FOR FINITE SEQUENCES

Let p be a finite sequence and let m, n be natural numbers. The functor $\langle p(m), \dots, p(n) \rangle$ yields a finite sequence and is defined by:

- (Def.1) (i) $\text{len}\langle p(m), \dots, p(n) \rangle + m = n + 1$ and for every natural number i such that $i < \text{len}\langle p(m), \dots, p(n) \rangle$ holds $\langle p(m), \dots, p(n) \rangle(i + 1) = p(m + i)$ if $1 \leq m$ and $m \leq n + 1$ and $n \leq \text{len } p$,
- (ii) $\langle p(m), \dots, p(n) \rangle = \varepsilon$, otherwise.

We now state several propositions:

- (6) If $1 \leq m$ and $m \leq \text{len } p$, then $\langle p(m), \dots, p(m) \rangle = \langle p(m) \rangle$.
- (7) $\langle p(1), \dots, p(\text{len } p) \rangle = p$.
- (8) If $m \leq n$ and $n \leq r$ and $r \leq \text{len } p$, then $\langle p(m + 1), \dots, p(n) \rangle \wedge \langle p(n + 1), \dots, p(r) \rangle = \langle p(m + 1), \dots, p(r) \rangle$.
- (9) If $1 \leq m$ and $m \leq \text{len } p$, then $\langle p(1), \dots, p(m) \rangle \wedge \langle p(m + 1), \dots, p(\text{len } p) \rangle = p$.
- (10) If $1 \leq m$ and $m \leq n$ and $n \leq \text{len } p$, then $\langle p(1), \dots, p(m) \rangle \wedge \langle p(m + 1), \dots, p(n) \rangle \wedge \langle p(n + 1), \dots, p(\text{len } p) \rangle = p$.
- (11) $\text{rng}\langle p(m), \dots, p(n) \rangle \subseteq \text{rng } p$.

Let D be a set, let p be a finite sequence of elements of D , and let m, n be natural numbers. Then $\langle p(m), \dots, p(n) \rangle$ is a finite sequence of elements of D .

Next we state the proposition

- (12) If $p \neq \varepsilon$ and $1 \leq m$ and $m \leq n$ and $n \leq \text{len } p$, then $\langle p(m), \dots, p(n) \rangle(1) = p(m)$ and $\langle p(m), \dots, p(n) \rangle(\text{len } \langle p(m), \dots, p(n) \rangle) = p(n)$.

3. THE GLUEING CATENATION OF FINITE SEQUENCES

Let p, q be finite sequences. The functor $p \curvearrowright q$ yielding a finite sequence is defined as follows:

(Def.2) $p \curvearrowright q = p \hat{\ } \langle q(2), \dots, q(\text{len } q) \rangle$.

Next we state several propositions:

- (13) If $q \neq \varepsilon$, then $\text{len}(p \curvearrowright q) + 1 = \text{len } p + \text{len } q$.
 (14) If $1 \leq k$ and $k \leq \text{len } p$, then $(p \curvearrowright q)(k) = p(k)$.
 (15) If $1 \leq k$ and $k < \text{len } q$, then $(p \curvearrowright q)(\text{len } p + k) = q(k + 1)$.
 (16) If $1 < \text{len } q$, then $(p \curvearrowright q)(\text{len}(p \curvearrowright q)) = q(\text{len } q)$.
 (17) $\text{rng}(p \curvearrowright q) \subseteq \text{rng } p \cup \text{rng } q$.

Let D be a set and let p, q be finite sequences of elements of D . Then $p \curvearrowright q$ is a finite sequence of elements of D .

Next we state the proposition

- (18) If $p \neq \varepsilon$ and $q \neq \varepsilon$ and $p(\text{len } p) = q(1)$, then $\text{rng}(p \curvearrowright q) = \text{rng } p \cup \text{rng } q$.

4. TWO VALUED ALTERNATING FINITE SEQUENCES

A finite sequence is two-valued if:

(Def.3) $\text{card } \text{rng } \text{it} = 2$.

The following proposition is true

- (19) p is two-valued iff $\text{len } p > 1$ and there exist arbitrary x, y such that $x \neq y$ and $\text{rng } p = \{x, y\}$.

A finite sequence is alternating if:

(Def.4) For every natural number i such that $1 \leq i$ and $i + 1 \leq \text{len } \text{it}$ it holds $\text{it}(i) \neq \text{it}(i + 1)$.

One can check that there exists a finite sequence which is two-valued and alternating.

In the sequel a, a_1, a_2 are two-valued alternating finite sequences.

One can prove the following propositions:

- (20) If $\text{len } a_1 = \text{len } a_2$ and $\text{rng } a_1 = \text{rng } a_2$ and $a_1(1) = a_2(1)$, then $a_1 = a_2$.
 (21) If $a_1 \neq a_2$ and $\text{len } a_1 = \text{len } a_2$ and $\text{rng } a_1 = \text{rng } a_2$, then for every i such that $1 \leq i$ and $i \leq \text{len } a_1$ holds $a_1(i) \neq a_2(i)$.
 (22) If $a_1 \neq a_2$ and $\text{len } a_1 = \text{len } a_2$ and $\text{rng } a_1 = \text{rng } a_2$, then for every a such that $\text{len } a = \text{len } a_1$ and $\text{rng } a = \text{rng } a_1$ holds $a = a_1$ or $a = a_2$.

- (23) If $X \neq Y$ and $n > 1$, then there exists a_1 such that $\text{rng } a_1 = \{X, Y\}$ and $\text{len } a_1 = n$ and $a_1(1) = X$.

5. FINITE SUBSEQUENCE OF FINITE SEQUENCES

Let us consider X and let f_1 be a finite sequence of elements of X . A finite subsequence is called a FinSubsequence of f_1 if:

(Def.5) It $\subseteq f_1$.

In the sequel s_1 will denote a finite subsequence.

The following propositions are true:

- (24) If s_1 is a finite sequence, then $\text{Seq } s_1 = s_1$.
 (25) If $\text{rng } p \subseteq \text{dom } s_1$, then $s_1 \cdot p$ is a finite sequence.
 (26) Let f be a finite subsequence and let g, h, f_2, f_3, f_4 be finite sequences. If $\text{rng } g \subseteq \text{dom } f$ and $\text{rng } h \subseteq \text{dom } f$ and $f_2 = f \cdot g$ and $f_3 = f \cdot h$ and $f_4 = f \cdot (g \wedge h)$, then $f_4 = f_2 \wedge f_3$.

We follow the rules: f_1, f_5, f_6 will be finite sequences of elements of X and f_7, f_8 will be FinSubsequence of f_1 .

We now state four propositions:

- (27) $\text{dom } f_7 \subseteq \text{dom } f_1$ and $\text{rng } f_7 \subseteq \text{rng } f_1$.
 (28) f_1 is a FinSubsequence of f_1 .
 (29) $f_7 \upharpoonright Y$ is a FinSubsequence of f_1 .
 (30) For every FinSubsequence f_9 of f_5 such that $\text{Seq } f_7 = f_5$ and $\text{Seq } f_9 = f_6$ and $f_8 = f_7 \upharpoonright \text{rng}(\text{Sgm dom } f_7 \upharpoonright \text{dom } f_9)$ holds $\text{Seq } f_8 = f_6$.

6. VERTEX SEQUENCES INDUCED BY CHAINS

In the sequel G is a graph.

Let us consider G . One can verify that the vertices of G is non empty.

We follow the rules: v, v_1, v_2, v_3, v_4 will denote elements of the vertices of G and e will be arbitrary.

We now state two propositions:

- (31) If e joins v_1 with v_2 , then e joins v_2 with v_1 .
 (32) If e joins v_1 with v_2 and e joins v_3 with v_4 , then $v_1 = v_3$ and $v_2 = v_4$ or $v_1 = v_4$ and $v_2 = v_3$.

Let us consider G . We see that the chain of G is a finite sequence of elements of the edges of G .

Let us consider G . A path of G is a path-like chain of G .

We follow the rules: v_5, v_6, v_7 will denote finite sequences of elements of the vertices of G and c, c_1, c_2 will denote chains of G .

The following proposition is true

(33) ε is a chain of G .

Let us consider G . One can check that there exists a chain of G which is empty.

Let us consider G, X . The functor $(G)\text{-VSet}(X)$ yields a set and is defined as follows:

(Def.6) $(G)\text{-VSet}(X) = \{v : \bigvee_{e: \text{element of the edges of } G} e \in X \wedge (v = (\text{the source of } G)(e) \vee v = (\text{the target of } G)(e))\}$.

Let us consider G, v_5 and let c be a finite sequence. We say that v_5 is vertex sequence of c if and only if:

(Def.7) $\text{len } v_5 = \text{len } c + 1$ and for every n such that $1 \leq n$ and $n \leq \text{len } c$ holds $c(n)$ joins $\pi_n v_5$ with $\pi_{n+1} v_5$.

One can prove the following four propositions:

(34) If $c \neq \varepsilon$ and v_5 is vertex sequence of c , then $(G)\text{-VSet}(\text{rng } c) = \text{rng } v_5$.

(35) $\langle v \rangle$ is vertex sequence of ε .

(36) There exists v_5 which is vertex sequence of c .

(37) Suppose $c \neq \varepsilon$ and v_6 is vertex sequence of c and v_7 is vertex sequence of c and $v_6 \neq v_7$. Then $v_6(1) \neq v_7(1)$ and for every v_5 such that v_5 is vertex sequence of c holds $v_5 = v_6$ or $v_5 = v_7$.

Let us consider G and let c be a finite sequence. We say that c alternates vertices in G if and only if:

(Def.8) $\text{len } c \geq 1$ and $\overline{\overline{(G)\text{-VSet}(\text{rng } c)}} = 2$ and for every n such that $n \in \text{dom } c$ holds $(\text{the source of } G)(c(n)) \neq (\text{the target of } G)(c(n))$.

One can prove the following propositions:

(38) If c alternates vertices in G and v_5 is vertex sequence of c , then for every k such that $k \in \text{dom } c$ holds $v_5(k) \neq v_5(k + 1)$.

(39) Suppose c alternates vertices in G and v_5 is vertex sequence of c . Then $\text{rng } v_5 = \{(\text{the source of } G)(c(1)), (\text{the target of } G)(c(1))\}$.

(40) Suppose c alternates vertices in G and v_5 is vertex sequence of c . Then v_5 is a two-valued alternating finite sequence.

(41) Suppose c alternates vertices in G . Then there exist v_6, v_7 such that

- (i) $v_6 \neq v_7$,
- (ii) v_6 is vertex sequence of c ,
- (iii) v_7 is vertex sequence of c , and
- (iv) for every v_5 such that v_5 is vertex sequence of c holds $v_5 = v_6$ or $v_5 = v_7$.

(42) Suppose v_5 is vertex sequence of c . Then $\overline{\overline{\text{the vertices of } G}} = 1$ or $c \neq \varepsilon$ and c does not alternate vertices in G if and only if for every v_6 such that v_6 is vertex sequence of c holds $v_6 = v_5$.

Let us consider G, c . Let us assume that $\overline{\overline{\text{the vertices of } G}} = 1$ or $c \neq \varepsilon$ and c does not alternate vertices in G . The functor $\text{vertex-seq}(c)$ yielding a finite sequence of elements of the vertices of G is defined as follows:

(Def.9) $\text{vertex-seq}(c)$ is vertex sequence of c .

We now state several propositions:

- (43) If v_5 is vertex sequence of c and $c_1 = c \upharpoonright \text{Seg } n$ and $v_6 = v_5 \upharpoonright \text{Seg}(n+1)$, then v_6 is vertex sequence of c_1 .
- (44) If $1 \leq m$ and $m \leq n$ and $n \leq \text{len } c$ and $q = \langle c(m), \dots, c(n) \rangle$, then q is a chain of G .
- (45) If $1 \leq m$ and $m \leq n$ and $n \leq \text{len } c$ and $c_1 = \langle c(m), \dots, c(n) \rangle$ and v_5 is vertex sequence of c and $v_6 = \langle v_5(m), \dots, v_5(n+1) \rangle$, then v_6 is vertex sequence of c_1 .
- (46) If v_6 is vertex sequence of c_1 and v_7 is vertex sequence of c_2 and $v_6(\text{len } v_6) = v_7(1)$, then $c_1 \hat{\ } c_2$ is a chain of G .
- (47) Suppose v_6 is vertex sequence of c_1 and v_7 is vertex sequence of c_2 and $v_6(\text{len } v_6) = v_7(1)$ and $c = c_1 \hat{\ } c_2$ and $v_5 = v_6 \curvearrowright v_7$. Then v_5 is vertex sequence of c .

7. VERTEX SEQUENCES INDUCED BY SIMPLE CHAINS, PATHS AND ORDERED CHAINS

Let us consider G . A chain of G is simple if it satisfies the condition (Def.10).

(Def.10) There exists v_5 such that v_5 is vertex sequence of it and for all n, m such that $1 \leq n$ and $n < m$ and $m \leq \text{len } v_5$ and $v_5(n) = v_5(m)$ holds $n = 1$ and $m = \text{len } v_5$.

Let us consider G . One can check that there exists a chain of G which is simple.

In the sequel s_2 denotes a simple chain of G .

Next we state several propositions:

- (49)² $s_2 \upharpoonright \text{Seg } n$ is a simple chain of G .
- (50) If $2 < \text{len } s_2$ and v_6 is vertex sequence of s_2 and v_7 is vertex sequence of s_2 , then $v_6 = v_7$.
- (51) If v_5 is vertex sequence of s_2 , then for all n, m such that $1 \leq n$ and $n < m$ and $m \leq \text{len } v_5$ and $v_5(n) = v_5(m)$ holds $n = 1$ and $m = \text{len } v_5$.
- (52) Suppose c is not a simple chain of G and v_5 is vertex sequence of c . Then there exists a FinSubsequence f_{10} of c and there exists a FinSubsequence f_{11} of v_5 and there exist c_1, v_6 such that $\text{len } c_1 < \text{len } c$ and v_6 is vertex sequence of c_1 and $\text{len } v_6 < \text{len } v_5$ and $v_5(1) = v_6(1)$ and $v_5(\text{len } v_5) = v_6(\text{len } v_6)$ and $\text{Seq } f_{10} = c_1$ and $\text{Seq } f_{11} = v_6$.
- (53) Suppose v_5 is vertex sequence of c . Then there exists a FinSubsequence f_{10} of c and there exists a FinSubsequence f_{11} of v_5 and there exist s_2, v_6 such that $\text{Seq } f_{10} = s_2$ and $\text{Seq } f_{11} = v_6$ and v_6 is vertex sequence of s_2 and $v_5(1) = v_6(1)$ and $v_5(\text{len } v_5) = v_6(\text{len } v_6)$.

²The proposition (48) has been removed.

Let us consider G . One can check that every chain of G which is empty is also path-like.

We now state the proposition

(54) If p is a path of G , then $p \upharpoonright \text{Seg } n$ is a path of G .

Let us consider G . One can verify that there exists a path of G which is simple.

We now state two propositions:

(55) If $2 < \text{len } s_2$, then s_2 is a path of G .

(56) s_2 is a path of G iff $\text{len } s_2 = 0$ or $\text{len } s_2 = 1$ or $s_2(1) \neq s_2(2)$.

Let us consider G . Observe that every chain of G which is empty is also oriented.

Let us consider G and let o_1 be an oriented chain of G . Let us assume that $o_1 \neq \varepsilon$. The functor $\text{vertex-seq}(o_1)$ yields a finite sequence of elements of the vertices of G and is defined as follows:

(Def.11) $\text{vertex-seq}(o_1)$ is vertex sequence of o_1 and $(\text{vertex-seq}(o_1))(1) = (\text{the source of } G)(o_1(1))$.

REFERENCES

- [1] Grzegorz Bancerek. Cardinal numbers. *Formalized Mathematics*, 1(2):377–382, 1990.
- [2] Grzegorz Bancerek. The fundamental properties of natural numbers. *Formalized Mathematics*, 1(1):41–46, 1990.
- [3] Grzegorz Bancerek. Zermelo theorem and axiom of choice. *Formalized Mathematics*, 1(2):265–267, 1990.
- [4] Grzegorz Bancerek and Krzysztof Hryniewiecki. Segments of natural numbers and finite sequences. *Formalized Mathematics*, 1(1):107–114, 1990.
- [5] Czesław Byliński. Finite sequences and tuples of elements of a non-empty sets. *Formalized Mathematics*, 1(3):529–536, 1990.
- [6] Czesław Byliński. Functions and their basic properties. *Formalized Mathematics*, 1(1):55–65, 1990.
- [7] Czesław Byliński. Functions from a set to a set. *Formalized Mathematics*, 1(1):153–164, 1990.
- [8] Czesław Byliński. Partial functions. *Formalized Mathematics*, 1(2):357–367, 1990.
- [9] Czesław Byliński. Some basic properties of sets. *Formalized Mathematics*, 1(1):47–53, 1990.
- [10] Agata Darmochwał. Finite sets. *Formalized Mathematics*, 1(1):165–167, 1990.
- [11] Krzysztof Hryniewiecki. Basic properties of real numbers. *Formalized Mathematics*, 1(1):35–40, 1990.
- [12] Krzysztof Hryniewiecki. Graphs. *Formalized Mathematics*, 2(3):365–370, 1991.
- [13] Takaya Nishiyama and Yasuho Mizuhara. Binary arithmetics. *Formalized Mathematics*, 4(1):83–86, 1993.
- [14] Andrzej Trybulec. Domains and their Cartesian products. *Formalized Mathematics*, 1(1):115–122, 1990.
- [15] Andrzej Trybulec. Enumerated sets. *Formalized Mathematics*, 1(1):25–34, 1990.
- [16] Andrzej Trybulec. Tarski Grothendieck set theory. *Formalized Mathematics*, 1(1):9–11, 1990.
- [17] Andrzej Trybulec. Tuples, projections and Cartesian products. *Formalized Mathematics*, 1(1):97–105, 1990.
- [18] Michał J. Trybulec. Integers. *Formalized Mathematics*, 1(3):501–505, 1990.
- [19] Wojciech A. Trybulec. Pigeon hole principle. *Formalized Mathematics*, 1(3):575–579, 1990.

- [20] Zinaida Trybulec and Halina Świączkowska. Boolean properties of sets. *Formalized Mathematics*, 1(1):17–23, 1990.
- [21] Edmund Woronowicz. Relations and their basic properties. *Formalized Mathematics*, 1(1):73–83, 1990.

Received May 13, 1995
