

Miscellaneous Facts about Functions

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The papers [16], [26], [3], [24], [29], [14], [28], [19], [23], [25], [22], [1], [17], [18], [30], [10], [6], [5], [15], [8], [13], [7], [11], [21], [9], [12], [2], [27], [20], and [4] provide the terminology and notation for this paper.

1. PRELIMINARIES

For simplicity we adopt the following rules: x is arbitrary, m, n are natural numbers, f, g are functions, and A, B are sets.

We now state several propositions:

- (1) For every function f and for every set X such that $\text{rng } f \subseteq X$ holds $\text{id}_X \cdot f = f$.
- (2) Let X be a set, and let Y be a non empty set, and let f be a function from X into Y . Suppose f is one-to-one. Let B be a subset of X and let C be a subset of Y . If $C \subseteq f^\circ B$, then $f^{-1} C \subseteq B$.
- (3) Let X, Y be non empty sets and let f be a function from X into Y . Suppose f is one-to-one. Let x be an element of X and let A be a subset of X . If $f(x) \in f^\circ A$, then $x \in A$.
- (4) Let X, Y be non empty sets and let f be a function from X into Y . Suppose f is one-to-one. Let x be an element of X , and let A be a subset of X , and let B be a subset of Y . If $f(x) \in f^\circ A \setminus B$, then $x \in A \setminus f^{-1} B$.
- (5) Let X, Y be non empty sets and let f be a function from X into Y . Suppose f is one-to-one. Let y be an element of Y , and let A be a subset of X , and let B be a subset of Y . If $y \in f^\circ A \setminus B$, then $f^{-1}(y) \in A \setminus f^{-1} B$.
- (6) For every function f and for arbitrary a such that $a \in \text{dom } f$ holds $f \upharpoonright \{a\} = a \mapsto f(a)$.

Let x, y be arbitrary. Observe that $x \dashrightarrow y$ is non empty.

Let x, y, a, b be arbitrary. One can check that $[x \dashrightarrow a, y \dashrightarrow b]$ is non empty.

One can prove the following propositions:

- (7) For every set I and for every many sorted set M indexed by I and for arbitrary i such that $i \in I$ holds $i \dashrightarrow M(i) = M \upharpoonright \{i\}$.
- (8) Let I, J be sets, and let M be a many sorted set indexed by $[I, J]$, and let i, j be arbitrary. If $i \in I$ and $j \in J$, then $[\langle i, j \rangle \dashrightarrow M(i, j)] = M \upharpoonright [\{i\}, \{j\}]$.
- (9) If $x \in \text{dom } f$ and $x \notin \text{dom } g$, then $(f + \cdot g)(x) = f(x)$.
- (10) For all functions f, g, h such that $\text{rng } g \subseteq \text{dom } f$ and $\text{rng } h \subseteq \text{dom } f$ holds $f \cdot (g + \cdot h) = f \cdot g + \cdot f \cdot h$.
- (11) For all functions f, g, h holds $(g + \cdot h) \cdot f = g \cdot f + \cdot h \cdot f$.
- (12) For all functions f, g, h such that $\text{rng } f$ misses $\text{dom } g$ holds $(h + \cdot g) \cdot f = h \cdot f$.
- (13) For all sets A, B and for arbitrary y such that A meets $\text{rng}(\text{id}_B + \cdot (A \dashrightarrow y))$ holds $y \in A$.
- (14) For arbitrary x, y and for every set A such that $x \neq y$ holds $x \notin \text{rng}(\text{id}_A + \cdot (x \dashrightarrow y))$.
- (15) For every set X and for arbitrary a and for every function f such that $\text{dom } f = X \cup \{a\}$ holds $f = f \upharpoonright X + \cdot (a \dashrightarrow f(a))$.
- (16) For every function f and for all sets X, y, z holds $f + \cdot (X \dashrightarrow y) + \cdot (X \dashrightarrow z) = f + \cdot (X \dashrightarrow z)$.
- (17) If $0 < m$ and $m \leq n$, then $\mathbb{Z}_m \subseteq \mathbb{Z}_n$.
- (18) $\mathbb{Z} \neq \mathbb{Z}^*$.
- (19) $\emptyset^* = \{\emptyset\}$.
- (20) $\langle x \rangle \in A^*$ iff $x \in A$.
- (21) $A \subseteq B$ iff $A^* \subseteq B^*$.
- (22) For every subset A of \mathbb{N} such that for all n, m such that $n \in A$ and $m < n$ holds $m \in A$ holds A is a cardinal number.
- (23) Let A be a finite set and let X be a non empty family of subsets of A . Then there exists an element C of X such that for every element B of X such that $B \subseteq C$ holds $B = C$.
- (24) Let p, q be finite sequences. Suppose $\text{len } p = \text{len } q + 1$. Let i be a natural number. Then $i \in \text{dom } q$ if and only if the following conditions are satisfied:
 - (i) $i \in \text{dom } p$, and
 - (ii) $i + 1 \in \text{dom } p$.

Let us note that there exists a finite sequence which is function yielding non empty and non-empty.

Note that ε is function yielding. Let f be a function. Observe that $\langle f \rangle$ is function yielding. Let g be a function. One can check that $\langle f, g \rangle$ is function

yielding. Let h be a function. Observe that $\langle f, g, h \rangle$ is function yielding.

Let n be a natural number and let f be a function. One can verify that $n \mapsto f$ is function yielding.

Let p be a finite sequence and let q be a non empty finite sequence. One can verify that $p \hat{\ } q$ is non empty and $q \hat{\ } p$ is non empty.

Let p, q be function yielding finite sequences. Note that $p \hat{\ } q$ is function yielding.

Next we state the proposition

- (25) Let p, q be finite sequences. Suppose $p \hat{\ } q$ is function yielding. Then p is function yielding and q is function yielding.

2. SOME USEFUL SCHEMES

In this article we present several logical schemes. The scheme *KappaD* concerns non empty sets \mathcal{A}, \mathcal{B} and a unary functor \mathcal{F} yielding arbitrary, and states that:

There exists a function f from \mathcal{A} into \mathcal{B} such that for every element x of \mathcal{A} holds $f(x) = \mathcal{F}(x)$

provided the parameters meet the following condition:

- For every element x of \mathcal{A} holds $\mathcal{F}(x) \in \mathcal{B}$.

The scheme *Kappa2D* deals with non empty sets $\mathcal{A}, \mathcal{B}, \mathcal{C}$ and a binary functor \mathcal{F} yielding arbitrary, and states that:

There exists a function f from $[\mathcal{A}, \mathcal{B}]$ into \mathcal{C} such that for every element x of \mathcal{A} and for every element y of \mathcal{B} holds $f(\langle x, y \rangle) = \mathcal{F}(x, y)$

provided the parameters meet the following requirement:

- For every element x of \mathcal{A} and for every element y of \mathcal{B} holds $\mathcal{F}(x, y) \in \mathcal{C}$.

The scheme *FinMono* concerns a set \mathcal{A} , a non empty set \mathcal{B} , and two unary functors \mathcal{F} and \mathcal{G} yielding arbitrary, and states that:

$\{\mathcal{F}(d) : d \text{ ranges over elements of } \mathcal{B}, \mathcal{G}(d) \in \mathcal{A}\}$ is finite

provided the following conditions are satisfied:

- \mathcal{A} is finite,
- For all elements d_1, d_2 of \mathcal{B} such that $\mathcal{G}(d_1) = \mathcal{G}(d_2)$ holds $d_1 = d_2$.

The scheme *CardMono* concerns a set \mathcal{A} , a non empty set \mathcal{B} , and a unary functor \mathcal{F} yielding arbitrary, and states that:

$\mathcal{A} \approx \{d : d \text{ ranges over elements of } \mathcal{B}, \mathcal{F}(d) \in \mathcal{A}\}$

provided the following requirements are met:

- For arbitrary x such that $x \in \mathcal{A}$ there exists an element d of \mathcal{B} such that $x = \mathcal{F}(d)$,
- For all elements d_1, d_2 of \mathcal{B} such that $\mathcal{F}(d_1) = \mathcal{F}(d_2)$ holds $d_1 = d_2$.

The scheme *CardMono'* concerns a set \mathcal{A} , a non empty set \mathcal{B} , and a unary functor \mathcal{F} yielding arbitrary, and states that:

$A \approx \{\mathcal{F}(d) : d \text{ ranges over elements of } \mathcal{B}, d \in \mathcal{A}\}$

provided the following conditions are satisfied:

- $\mathcal{A} \subseteq \mathcal{B}$,
- For all elements d_1, d_2 of \mathcal{B} such that $\mathcal{F}(d_1) = \mathcal{F}(d_2)$ holds $d_1 = d_2$.

The scheme *FuncSeqInd* concerns a unary predicate \mathcal{P} , and states that:

For every function yielding finite sequence p holds $\mathcal{P}[p]$

provided the following conditions are satisfied:

- $\mathcal{P}[\varepsilon]$,
- For every function yielding finite sequence p such that $\mathcal{P}[p]$ and for every function f holds $\mathcal{P}[p \hat{\ } \langle f \rangle]$.

3. SOME AUXILIARY CONCEPTS

Let x be arbitrary and let y be a set. Let us assume that $x \in y$. The functor $x(\in y)$ yielding an element of y is defined as follows:

(Def. 1) $x(\in y) = x$.

One can prove the following proposition

(26) If $x \in A \cap B$, then $x(\in A) = x(\in B)$.

Let f, g be functions and let A be a set. We say that f and g equal outside A if and only if:

(Def. 2) $f \upharpoonright (\text{dom } f \setminus A) = g \upharpoonright (\text{dom } g \setminus A)$.

Next we state several propositions:

(27) For every function f and for every set A holds f and f equal outside A .

(28) For all functions f, g and for every set A such that f and g equal outside A holds g and f equal outside A

(29) Let f, g, h be functions and let A be a set. Suppose f and g equal outside A and g and h equal outside A . Then f and h equal outside A .

(30) For all functions f, g and for every set A such that f and g equal outside A holds $\text{dom } f \setminus A = \text{dom } g \setminus A$.

(31) For all functions f, g and for every set A such that $\text{dom } g \subseteq A$ holds f and $f + \cdot g$ equal outside A

Let f be a function and let i, x be arbitrary. The functor $f + \cdot (i, x)$ yields a function and is defined by:

(Def. 3) (i) $f + \cdot (i, x) = f + \cdot (i \mapsto x)$ if $i \in \text{dom } f$,
(ii) $f + \cdot (i, x) = f$, otherwise.

Next we state several propositions:

(32) For every function f and for arbitrary d, i holds $\text{dom}(f + \cdot (i, d)) = \text{dom } f$.

(33) For every function f and for arbitrary d, i such that $i \in \text{dom } f$ holds $(f + \cdot (i, d))(i) = d$.

- (34) For every function f and for arbitrary d, i, j such that $i \neq j$ and $j \in \text{dom } f$ holds $(f + \cdot (i, d))(j) = f(j)$.
- (35) For every function f and for arbitrary d, e, i, j such that $i \neq j$ holds $f + \cdot (i, d) + \cdot (j, e) = f + \cdot (j, e) + \cdot (i, d)$.
- (36) For every function f and for arbitrary d, e, i holds $f + \cdot (i, d) + \cdot (i, e) = f + \cdot (i, e)$.
- (37) For every function f and for arbitrary i holds $f + \cdot (i, f(i)) = f$.

Let f be a finite sequence, let i be a natural number, and let x be arbitrary. One can check that $f + \cdot (i, x)$ is finite sequence-like.

Let D be a set, let f be a finite sequence of elements of D , let i be a natural number, and let d be an element of D . Then $f + \cdot (i, d)$ is a finite sequence of elements of D .

The following three propositions are true:

- (38) Let D be a non empty set, and let f be a finite sequence of elements of D , and let d be an element of D , and let i be a natural number. If $i \in \text{dom } f$, then $\pi_i(f + \cdot (i, d)) = d$.
- (39) Let D be a non empty set, and let f be a finite sequence of elements of D , and let d be an element of D , and let i, j be natural numbers. If $i \neq j$ and $j \in \text{dom } f$, then $\pi_j(f + \cdot (i, d)) = \pi_j f$.
- (40) Let D be a non empty set, and let f be a finite sequence of elements of D , and let d, e be elements of D , and let i be a natural number. Then $f + \cdot (i, \pi_i f) = f$.

4. ON THE COMPOSITION OF A FINITE SEQUENCE OF FUNCTIONS

Let X be a set and let p be a function yielding finite sequence. The functor $\text{compose}_X p$ yielding a function is defined by the condition (Def. 4).

(Def. 4) There exists a many sorted function f of \mathbb{N} such that

- (i) $\text{compose}_X p = f(\text{len } p)$,
- (ii) $f(0) = \text{id}_X$, and
- (iii) for every natural number i such that $i + 1 \in \text{dom } p$ and for all functions g, h such that $g = f(i)$ and $h = p(i + 1)$ holds $f(i + 1) = h \cdot g$.

Let p be a function yielding finite sequence and let x be a set. The functor $\text{apply}(p, x)$ yields a finite sequence and is defined by the conditions (Def. 5).

- (Def. 5) (i) $\text{len } \text{apply}(p, x) = \text{len } p + 1$,
- (ii) $(\text{apply}(p, x))(1) = x$, and
 - (iii) for every natural number i and for every function f such that $i \in \text{dom } p$ and $f = p(i)$ holds $(\text{apply}(p, x))(i + 1) = f((\text{apply}(p, x))(i))$.

We adopt the following convention: X, Y, x denote sets, p, q denote function yielding finite sequences, and f, g, h denote functions.

The following propositions are true:

- (41) $\text{compose}_X \varepsilon = \text{id}_X$.
- (42) $\text{apply}(\varepsilon, x) = \langle x \rangle$.
- (43) $\text{compose}_X(p \wedge \langle f \rangle) = f \cdot \text{compose}_X p$.
- (44) $\text{apply}(p \wedge \langle f \rangle, x) = (\text{apply}(p, x)) \wedge \langle f((\text{apply}(p, x))(\text{len } p + 1)) \rangle$.
- (45) $\text{compose}_X(\langle f \rangle \wedge p) = \text{compose}_{f \circ X} p \cdot (f \upharpoonright X)$.
- (46) $\text{apply}(\langle f \rangle \wedge p, x) = \langle x \rangle \wedge \text{apply}(p, f(x))$.
- (47) $\text{compose}_X \langle f \rangle = f \cdot \text{id}_X$.
- (48) If $\text{dom } f \subseteq X$, then $\text{compose}_X \langle f \rangle = f$.
- (49) $\text{apply}(\langle f \rangle, x) = \langle x, f(x) \rangle$.
- (50) If $\text{rng } \text{compose}_X p \subseteq Y$, then $\text{compose}_X(p \wedge q) = \text{compose}_Y q \cdot \text{compose}_X p$.
- (51) $(\text{apply}(p \wedge q, x))(\text{len}(p \wedge q) + 1) = (\text{apply}(q, (\text{apply}(p, x))(\text{len } p + 1)))(\text{len } q + 1)$.
- (52) $\text{apply}(p \wedge q, x) = (\text{apply}(p, x))^{\S \wedge} \text{apply}(q, (\text{apply}(p, x))(\text{len } p + 1))$.
- (53) $\text{compose}_X \langle f, g \rangle = g \cdot f \cdot \text{id}_X$.
- (54) If $\text{dom } f \subseteq X$ or $\text{dom}(g \cdot f) \subseteq X$, then $\text{compose}_X \langle f, g \rangle = g \cdot f$.
- (55) $\text{apply}(\langle f, g \rangle, x) = \langle x, f(x), g(f(x)) \rangle$.
- (56) $\text{compose}_X \langle f, g, h \rangle = h \cdot g \cdot f \cdot \text{id}_X$.
- (57) If $\text{dom } f \subseteq X$ or $\text{dom}(g \cdot f) \subseteq X$ or $\text{dom}(h \cdot g \cdot f) \subseteq X$, then $\text{compose}_X \langle f, g, h \rangle = h \cdot g \cdot f$.
- (58) $\text{apply}(\langle f, g, h \rangle, x) = \langle x \rangle \wedge \langle f(x), g(f(x)), h(g(f(x))) \rangle$.

Let F be a finite sequence. The functor $\text{firstdom}(F)$ is defined as follows:

- (Def. 6) (i) $\text{firstdom}(F)$ is empty if F is empty,
(ii) $\text{firstdom}(F) = \pi_1(F(1))$, otherwise.

The functor $\text{lastrng}(F)$ is defined by:

- (Def. 7) (i) $\text{lastrng}(F)$ is empty if F is empty,
(ii) $\text{lastrng}(F) = \pi_2(F(\text{len } F))$, otherwise.

Next we state three propositions:

- (59) $\text{firstdom}(\varepsilon) = \emptyset$ and $\text{lastrng}(\varepsilon) = \emptyset$.
- (60) For every finite sequence p holds $\text{firstdom}(\langle f \rangle \wedge p) = \text{dom } f$ and $\text{lastrng}(p \wedge \langle f \rangle) = \text{rng } f$.
- (61) For every function yielding finite sequence p such that $p \neq \varepsilon$ holds $\text{rng } \text{compose}_X p \subseteq \text{lastrng}(p)$.

Let I_1 be a finite sequence. We say that I_1 is composable if and only if:

- (Def. 8) There exists a finite sequence p such that $\text{len } p = \text{len } I_1 + 1$ and for every natural number i such that $i \in \text{dom } I_1$ holds $I_1(i) \in p(i + 1)^{p(i)}$.

We now state the proposition

- (62) For all finite sequences p, q such that $p \wedge q$ is composable holds p is composable and q is composable.

One can verify that every finite sequence which is composable is also function yielding.

Let us observe that every finite sequence which is empty is also composable.

Let f be a function. One can check that $\langle f \rangle$ is composable.

Let us observe that there exists a finite sequence which is composable non empty and non-empty.

A composable sequence is a composable finite sequence.

Next we state several propositions:

- (63) For every composable sequence p such that $p \neq \varepsilon$ holds $\text{dom compose}_X p = \text{firstdom}(p) \cap X$.
- (64) For every composable sequence p holds $\text{dom compose}_{\text{firstdom}(p)} p = \text{firstdom}(p)$.
- (65) For every composable sequence p and for every function f such that $\text{rng } f \subseteq \text{firstdom}(p)$ holds $\langle f \rangle \wedge p$ is a composable sequence.
- (66) For every composable sequence p and for every function f such that $\text{lastrng}(p) \subseteq \text{dom } f$ holds $p \wedge \langle f \rangle$ is a composable sequence.
- (67) For every composable sequence p such that $x \in \text{firstdom}(p)$ and $x \in X$ holds $(\text{apply}(p, x))(\text{len } p + 1) = (\text{compose}_X p)(x)$.

Let X, Y be sets. Let us assume that if Y is empty, then X is empty. A composable sequence is called a composable sequence from X into Y if:

(Def. 9) $\text{firstdom}(\text{it}) = X$ and $\text{lastrng}(\text{it}) \subseteq Y$.

Let Y be a non empty set, let X be a set, and let F be a composable sequence from X into Y . Then $\text{compose}_X F$ is a function from X into Y .

Let q be a non-empty non empty finite sequence. A finite sequence is said to be a composable sequence along q if:

(Def. 10) $\text{len it} + 1 = \text{len } q$ and for every natural number i such that $i \in \text{dom it}$ holds $\text{it}(i) \in q(i + 1)^{q(i)}$.

Let q be a non-empty non empty finite sequence. Observe that every composable sequence along q is composable and non-empty.

One can prove the following three propositions:

- (68) Let q be a non-empty non empty finite sequence and let p be a composable sequence along q . If $p \neq \varepsilon$, then $\text{firstdom}(p) = q(1)$ and $\text{lastrng}(p) \subseteq q(\text{len } q)$.
- (69) Let q be a non-empty non empty finite sequence and let p be a composable sequence along q . Then $\text{dom compose}_{q(1)} p = q(1)$ and $\text{rng compose}_{q(1)} p \subseteq q(\text{len } q)$.
- (70) For every function f and for every natural number n holds $f^n = \text{compose}_{\text{dom } f \cup \text{rng } f}(n \mapsto f)$.

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