Introduction to Concept Lattices

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Summary. In this paper we give Mizar formalization of concept lattices. Concept lattices stem from the so called formal concept analysis — a part of applied mathematics that brings mathematical methods into the field of data anylysis and knowledge processing. Our approach follows the one given in [8].

 ${\rm MML} \ {\rm Identifier:} \ {\tt CONLAT_1}.$

The papers [3], [14], [4], [5], [1], [15], [12], [10], [13], [11], [2], [7], [9], and [6] provide the notation and terminology for this paper.

1. Formal Contexts

We consider 2-sorted as systems

 $\langle \text{ objects, a Attributes } \rangle$,

where the objects constitute a set and the Attributes is a set.

Let C be a 2-sorted. We say that C is empty if and only if:

(Def. 1) The objects of C are empty and the Attributes of C is empty.

We say that C is quasi-empty if and only if:

(Def. 2) The objects of C are empty or the Attributes of C is empty.

Let us note that there exists a 2-sorted which is strict and non empty and there exists a 2-sorted which is strict and non quasi-empty.

One can verify that there exists a 2-sorted which is strict, empty, and quasiempty.

We consider ContextStr as extensions of 2-sorted as systems

 $\langle \text{ objects, a Attributes, a Information} \rangle$,

where the objects constitute a set, the Attributes is a set, and the Information is a relation between the objects and the Attributes.

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One can check that there exists a ContextStr which is strict and non empty and there exists a ContextStr which is strict and non quasi-empty.

A FormalContext is a non quasi-empty ContextStr.

Let C be a 2-sorted.

(Def. 3) An element of the objects of C is said to be an object of C.

(Def. 4) An element of the Attributes of C is said to be a Attribute of C.

Let C be a non quasi-empty 2-sorted. Note that the Attributes of C is non empty and the objects of C is non empty.

Let C be a non quasi-empty 2-sorted. One can check that there exists a subset of the objects of C which is non empty and there exists a subset of the Attributes of C which is non empty.

Let C be a FormalContext, let o be an object of C, and let a be a Attribute of C. We say that o is connected with a if and only if:

(Def. 5) $\langle o, a \rangle \in$ the Information of C.

We introduce o is not connected with a as an antonym of o is connected with a.

2. DERIVATION OPERATORS

Let C be a FormalContext. The functor ObjectDerivation C yields a function from $2^{\text{the objects of } C}$ into $2^{\text{the Attributes of } C}$ and is defined by the condition (Def. 6).

(Def. 6) Let O be an element of $2^{\text{the objects of } C}$. Then (ObjectDerivation C)(O) = $\{a; a \text{ ranges over Attribute of } C: \bigwedge_{o: \text{object of } C} (o \in O \Rightarrow o \text{ is connected with } a)\}.$

Let C be a FormalContext. The functor AttributeDerivation C yields a function from $2^{\text{the Attributes of } C}$ into $2^{\text{the objects of } C}$ and is defined by the condition (Def. 7).

(Def. 7) Let A be an element of 2^{the Attributes of C}. Then (AttributeDerivation C)(A) = $\{o; o \text{ ranges over objects of } C: \bigwedge_{a: \text{Attribute of } C} (a \in A \Rightarrow o \text{ is connected} with a)\}.$

The following propositions are true:

- (1) Let C be a FormalContext and o be an object of C. Then $(\text{ObjectDerivation } C)(\{o\}) = \{a; a \text{ ranges over Attribute of } C: o \text{ is connected with } a\}.$
- (2) Let C be a FormalContext and a be a Attribute of C. Then $(AttributeDerivation C)(\{a\}) = \{o; o \text{ ranges over objects of } C: o \text{ is connected with } a\}.$

- (3) For every FormalContext C and for all subsets O_1 , O_2 of the objects of C such that $O_1 \subseteq O_2$ holds $(\text{ObjectDerivation } C)(O_2) \subseteq (\text{ObjectDerivation } C)(O_1).$
- (4) For every FormalContext C and for all subsets A_1 , A_2 of the Attributes of C such that $A_1 \subseteq A_2$ holds (AttributeDerivation C) $(A_2) \subseteq$ (AttributeDerivation C) (A_1) .
- (5) For every FormalContext C and for every subset O of the objects of C holds $O \subseteq (AttributeDerivation <math>C)((ObjectDerivation C)(O)).$
- (6) For every FormalContext C and for every subset A of the Attributes of C holds $A \subseteq (\text{ObjectDerivation } C)((\text{AttributeDerivation } C)(A)).$
- (7) For every FormalContext C and for every subset O of the objects of C holds (ObjectDerivation C)(O) = (ObjectDerivation C) ((AttributeDerivation C)((ObjectDerivation C)(O))).
- (8) For every FormalContext C and for every subset A of the Attributes of C holds (AttributeDerivation C)(A) =
 (AttributeDerivation C)((ObjectDerivation C)((AttributeDerivation C)(A))).
- (9) Let C be a FormalContext, O be a subset of the objects of C, and A be a subset of the Attributes of C. Then $O \subseteq (\text{AttributeDerivation } C)(A)$ if and only if $[O, A] \subseteq$ the Information of C.
- (10) Let C be a FormalContext, O be a subset of the objects of C, and A be a subset of the Attributes of C. Then $A \subseteq (\text{ObjectDerivation } C)(O)$ if and only if $[O, A] \subseteq$ the Information of C.
- (11) Let C be a FormalContext, O be a subset of the objects of C, and A be a subset of the Attributes of C. Then $O \subseteq (\text{AttributeDerivation } C)(A)$ if and only if $A \subseteq (\text{ObjectDerivation } C)(O)$.

Let C be a Formal Context. The functor $\phi(C)$ yielding a map from $2_{\subset}^{\text{the objects of } C}$ into $2_{\subseteq}^{\text{the Attributes of } C}$ is defined by:

(Def. 8) $\phi(C) = \text{ObjectDerivation } C.$

Let C be a FormalContext. The functor psiC yields a map from $2_{\subseteq}^{\text{the Attributes of } C}$ into $2_{\subseteq}^{\text{the objects of } C}$ and is defined as follows:

(Def. 9) psi C = AttributeDerivation C.

We now state the proposition

(12) For every FormalContext C holds $\langle \phi(C), \operatorname{psi} C \rangle$ is a connection between $2_{\subset}^{\operatorname{the objects of } C}$ and $2_{\subseteq}^{\operatorname{the Attributes of } C}$.

Let P, R be non empty relational structures and let C_1 be a connection between P and R. We say that C_1 is co-Galois if and only if the condition (Def. 10) is satisfied.

(Def. 10) There exists a map f from P into R and there exists a map g from R into P such that

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- (i) $C_1 = \langle f, g \rangle$,
- (ii) f is antitone,
- (iii) g is antitone, and
- (iv) for all elements p_1 , p_2 of P and for all elements r_1 , r_2 of R holds $p_1 \leq g(f(p_1))$ and $r_1 \leq f(g(r_1))$.

We now state several propositions:

- (13) Let P, R be non empty posets, C_1 be a connection between P and R, f be a map from P into R, and g be a map from R into P. Suppose $C_1 = \langle f, g \rangle$. Then C_1 is co-Galois if and only if for every element p of P and for every element r of R holds $p \leq g(r)$ iff $r \leq f(p)$.
- (14) Let P, R be non empty posets and C_1 be a connection between P and R. Suppose C_1 is co-Galois. Let f be a map from P into R and g be a map from R into P. If $C_1 = \langle f, g \rangle$, then $f = f \cdot (g \cdot f)$ and $g = g \cdot (f \cdot g)$.
- (15) For every FormalContext C holds $\langle \phi(C), \operatorname{psi} C \rangle$ is co-Galois.
- (16) For every FormalContext C and for all subsets O_1 , O_2 of the objects of C holds (ObjectDerivation C) $(O_1 \cup O_2) = (ObjectDerivation <math>C)(O_1) \cap (ObjectDerivation C)(O_2)$.
- (17) For every FormalContext C and for all subsets A_1 , A_2 of the Attributes of C holds (AttributeDerivation C) $(A_1 \cup A_2) =$ (AttributeDerivation C) $(A_1) \cap$ (AttributeDerivation C) (A_2) .
- (18) For every FormalContext C holds (ObjectDerivation C)(\emptyset) = the Attributes of C.
- (19) For every FormalContext C holds (AttributeDerivation C)(\emptyset) = the objects of C.

3. Formal Concepts

Let C be a 2-sorted. We introduce ConceptStr over C which are systems $\langle a \text{ Extent}, a \text{ Intent} \rangle$,

where the Extent is a subset of the objects of C and the Intent is a subset of the Attributes of C.

Let C be a 2-sorted and let C_2 be a ConceptStr over C. We say that C_2 is empty if and only if:

(Def. 11) The Extent of C_2 is empty and the Intent of C_2 is empty.

We say that C_2 is quasi-empty if and only if:

(Def. 12) The Extent of C_2 is empty or the Intent of C_2 is empty.

Let C be a non quasi-empty 2-sorted. Observe that there exists a ConceptStr over C which is strict and non empty and there exists a ConceptStr over C which is strict and quasi-empty.

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Let C be an empty 2-sorted. Observe that every ConceptStr over C is empty. Let C be a quasi-empty 2-sorted. Observe that every ConceptStr over C is quasi-empty.

Let C be a FormalContext and let C_2 be a ConceptStr over C. We say that C_2 is concept-like if and only if:

(Def. 13) (ObjectDerivation C)(the Extent of C_2) = the Intent of C_2 and (AttributeDerivation C)(the Intent of C_2) = the Extent of C_2 .

Let C be a FormalContext. One can check that there exists a ConceptStr over C which is concept-like and non empty.

Let C be a FormalContext. A FormalConcept of C is a concept-like non empty ConceptStr over C.

Let C be a FormalContext. Note that there exists a FormalConcept of C which is strict.

Next we state four propositions:

- (20) Let C be a FormalContext and O be a subset of the objects of C. Then (i) $\langle (AttributeDerivation C)((ObjectDerivation C)(O)), \rangle$
 - (ObjectDerivation C)(O) is a FormalConcept of C, and
 - (ii) for every subset O' of the objects of C and for every subset A' of the Attributes of C such that $\langle O', A' \rangle$ is a FormalConcept of C and $O \subseteq O'$ holds (AttributeDerivation C)((ObjectDerivation C)(O)) $\subseteq O'$.
- (21) Let C be a FormalContext and O be a subset of the objects of C. Then there exists a subset A of the Attributes of C such that $\langle O, A \rangle$ is a FormalConcept of C if and only if (AttributeDerivation C)((ObjectDerivation C)(O)) = O.
- (22) Let C be a FormalContext and A be a subset of the Attributes of C. Then
 - (i) $\langle (AttributeDerivation C)(A), (ObjectDerivation C) \rangle ((AttributeDerivation C)(A)) \rangle$ is a FormalConcept of C, and
 - (ii) for every subset O' of the objects of C and for every subset A' of the Attributes of C such that $\langle O', A' \rangle$ is a FormalConcept of C and $A \subseteq A'$ holds (ObjectDerivation C)((AttributeDerivation C)(A) $\subseteq A'$.
- (23) Let C be a FormalContext and A be a subset of the Attributes of C. Then there exists a subset O of the objects of C such that $\langle O, A \rangle$ is a FormalConcept of C if and only if (ObjectDerivation C)((AttributeDerivation C)(A)) = A.

Let C be a FormalContext and let C_2 be a ConceptStr over C. We say that C_2 is universal if and only if:

(Def. 14) The Extent of C_2 = the objects of C.

Let C be a FormalContext and let C_2 be a ConceptStr over C. We say that C_2 is co-universal if and only if:

(Def. 15) The Intent of C_2 = the Attributes of C.

Let C be a FormalContext. Note that there exists a FormalConcept of C which is strict and universal and there exists a FormalConcept of C which is strict and co-universal.

Let C be a FormalContext. The functor Concept - with - all - Objects C yields a strict universal FormalConcept of C and is defined by the condition (Def. 16).

(Def. 16) There exists a subset O of the objects of C and there exists a subset A of the Attributes of C such that Concept – with – all – Objects $C = \langle O, A \rangle$ and $O = (AttributeDerivation <math>C)(\emptyset)$ and A =

 $(\text{ObjectDerivation}\,C)((\text{AttributeDerivation}\,C)(\emptyset)).$

Let C be a FormalContext. The functor Concept – with – all – Attributes C yielding a strict co-universal FormalConcept of C is defined by the condition (Def. 17).

(Def. 17) There exists a subset O of the objects of C and there exists a subset A of the Attributes of C such that Concept – with – all – Attributes $C = \langle O, A \rangle$ and $O = (AttributeDerivation <math>C)((ObjectDerivation C)(\emptyset))$ and $A = (ObjectDerivation C)(\emptyset)$.

One can prove the following propositions:

- (24) Let C be a FormalContext. Then the Extent of Concept – with – all – Objects C = the objects of C and the Intent of Concept – with – all – Attributes C = the Attributes of C.
- (25) Let C be a FormalContext and C_2 be a FormalConcept of C. Then
- (i) if the Extent of $C_2 = \emptyset$, then C_2 is co-universal, and
- (ii) if the Intent of $C_2 = \emptyset$, then C_2 is universal.
- (26) Let C be a FormalContext and C_2 be a strict FormalConcept of C. Then
 - (i) if the Extent of $C_2 = \emptyset$, then $C_2 = \text{Concept} \text{with} \text{all} \text{Attributes } C$, and
 - (ii) if the Intent of $C_2 = \emptyset$, then $C_2 = \text{Concept} \text{with} \text{all} \text{Objects } C$.
- (27) Let C be a FormalContext and C_2 be a quasi-empty ConceptStr over C. Suppose C_2 is a FormalConcept of C. Then C_2 is universal or co-universal.
- (28) Let C be a FormalContext and C_2 be a quasi-empty ConceptStr over C. If C_2 is a strict FormalConcept of C, then $C_2 = Concept with all Objects <math>C$ or $C_2 = Concept with all Attributes <math>C$.

Let C be a FormalContext. A non empty set is called a Set of FormalConcepts of C if:

(Def. 18) For every set X such that $X \in \text{it holds } X$ is a FormalConcept of C.

Let C be a FormalContext and let F_1 be a Set of FormalConcepts of C. We see that the element of F_1 is a FormalConcept of C.

Let C be a FormalContext and let C_3 , C_4 be FormalConcept of C. We say that C_3 is SubConcept of C_4 if and only if:

(Def. 19) The Extent of $C_3 \subseteq$ the Extent of C_4 .

We introduce C_4 is SuperConcept of C_3 as a synonym of C_3 is SubConcept of C_4 .

One can prove the following propositions:

- (29) Let C be a FormalContext and C_3 , C_4 be FormalConcept of C. Then C_3 is SuperConcept of C_4 if and only if C_4 is SubConcept of C_3 .
- (30) Let C be a FormalContext and C_3 , C_4 be FormalConcept of C. Then C_3 is SubConcept of C_4 if and only if the Extent of $C_3 \subseteq$ the Extent of C_4 .
- (31) Let C be a FormalContext and C_3 , C_4 be FormalConcept of C. Then C_3 is SubConcept of C_4 if and only if the Intent of $C_4 \subseteq$ the Intent of C_3 .
- (32) Let C be a FormalContext and C_3 , C_4 be FormalConcept of C. Then C_3 is SuperConcept of C_4 if and only if the Extent of $C_4 \subseteq$ the Extent of C_3 .
- (33) Let C be a FormalContext and C_3 , C_4 be FormalConcept of C. Then C_3 is SuperConcept of C_4 if and only if the Intent of $C_3 \subseteq$ the Intent of C_4 .
- (34) Let C be a FormalContext and C_2 be a FormalConcept of C. Then C_2 is SubConcept of Concept with all Objects C and Concept with all Attributes C is SubConcept of C_2 .

4. Concept Lattices

Let C be a FormalContext. The functor $B - \operatorname{carrier} C$ yielding a non empty set is defined by the condition (Def. 20).

- (Def. 20) B carrier $C = \{\langle E, I \rangle; E \text{ ranges over subsets of the objects of } C, I \text{ ranges over subsets of the Attributes of } C: \langle E, I \rangle \text{ is non empty } \land (\text{ObjectDerivation } C)(E) = I \land (\text{AttributeDerivation } C)(I) = E\}.$
 - Let C be a FormalContext. Then $B \operatorname{carrier} C$ is a Set of FormalConcepts of C.

Let C be a FormalContext. One can check that $B - \operatorname{carrier} C$ is non empty. One can prove the following proposition

(35) For every FormalContext C and for every set C_2 holds $C_2 \in B - \text{carrier } C$ iff C_2 is a strict FormalConcept of C.

Let C be a FormalContext. The functor B - meet C yields a binary operation on B - carrier C and is defined by the condition (Def. 21).

(Def. 21) Let C_3 , C_4 be strict FormalConcept of C. Then there exists a subset O of the objects of C and there exists a subset A of the Attributes of C such that

 $(B - meet C)(C_3, C_4) = \langle O, A \rangle$ and $O = (the Extent of C_3) \cap (the Extent of C_4)$ and $A = (ObjectDerivation C)((AttributeDerivation C))((the Intent of C_3) \cup (the Intent of C_4))).$

Let C be a FormalContext. The functor B - join C yielding a binary operation on B - carrier C is defined by the condition (Def. 22).

(Def. 22) Let C_3 , C_4 be strict FormalConcept of C. Then there exists a subset O of the objects of C and there exists a subset A of the Attributes of C such that $(B - join C)(C_3, C_4) = \langle O, A \rangle$ and O =(AttributeDerivation C)((ObjectDerivation C)((the Extent of C_3) \cup (the Extent of C_4))) and A = (the Intent of C_3) \cap (the Intent of C_4).

One can prove the following propositions:

- (36) For every FormalContext C and for all strict FormalConcept C_3 , C_4 of C holds $(B meet C)(C_3, C_4) = (B meet C)(C_4, C_3).$
- (37) For every FormalContext C and for all strict FormalConcept C_3 , C_4 of C holds $(B join C)(C_3, C_4) = (B join C)(C_4, C_3)$.
- (38) For every FormalContext C and for all strict FormalConcept C_3 , C_4 , C_5 of C holds $(B meet C)(C_3, (B meet C)(C_4, C_5)) = (B meet C)((B meet C)(C_3, C_4), C_5).$
- (39) For every FormalContext C and for all strict FormalConcept C_3 , C_4 , C_5 of C holds $(B join C)(C_3, (B join C)(C_4, C_5)) = (B join C)((B join C)(C_3, C_4), C_5).$
- (40) For every FormalContext C and for all strict FormalConcept C_3 , C_4 of C holds $(B join C)((B meet C)(C_3, C_4), C_4) = C_4$.
- (41) For every FormalContext C and for all strict FormalConcept C_3 , C_4 of C holds $(B meet C)(C_3, (B join C)(C_3, C_4)) = C_3$.
- (42) For every FormalContext C and for every strict FormalConcept C_2 of C holds $(B meet C)(C_2, Concept with all Objects C) = C_2.$
- (43) For every FormalContext C and for every strict FormalConcept C_2 of C holds $(B join C)(C_2, Concept with all Objects <math>C) = Concept with all Objects C.$
- (44) For every FormalContext C and for every strict FormalConcept C_2 of C holds $(B join C)(C_2, Concept with all Attributes C) = C_2$.
- (45) For every FormalContext C and for every strict FormalConcept C_2 of C holds $(B meet C)(C_2, Concept with all Attributes <math>C) = Concept with all Attributes <math>C$.

Let C be a FormalContext. The functor ConceptLattice C yielding a strict non empty lattice structure is defined as follows:

(Def. 23) ConceptLattice $C = \langle B - carrier C, B - join C, B - meet C \rangle$.

The following proposition is true

(46) For every FormalContext C holds ConceptLattice C is a lattice.

Let C be a FormalContext. One can verify that ConceptLattice C is strict non empty and lattice-like.

Let C be a FormalContext and let S be a non empty subset of the carrier of ConceptLattice C. We see that the element of S is a FormalConcept of C.

Let C be a FormalContext and let C_2 be an element of the carrier of ConceptLattice C. The functor C_2^{T} yielding a strict FormalConcept of C is defined as follows:

(Def. 24) $C_2^{\mathrm{T}} = C_2$.

One can prove the following two propositions:

- (47) Let C be a FormalContext and C_3 , C_4 be elements of the carrier of ConceptLattice C. Then $C_3 \sqsubseteq C_4$ if and only if C_3^{T} is SubConcept of C_4^{T} .
- (48) For every FormalContext C holds ConceptLattice C is a complete lattice.

Let C be a FormalContext. Observe that ConceptLattice C is complete.

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