

Basic Formal Properties of Triangular Norms and Conorms

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Summary. In the article we present in the Mizar system [1], [8] the catalogue of triangular norms and conorms, used especially in the theory of fuzzy sets [13]. The name *triangular* emphasizes the fact that in the framework of probabilistic metric spaces they generalize triangle inequality [2].

After defining corresponding Mizar mode using four attributes, we introduced the following t-norms:

- minimum t-norm minnorm (Def. 6),
- product t-norm prodnorm (Def. 8),
- Łukasiewicz t-norm Lukasiewicz_norm (Def. 10),
- drastic t-norm drastic_norm (Def. 11),
- nilpotent minimum nilmin_norm (Def. 12),
- Hamacher product Hamacher_norm (Def. 13),

and corresponding t-conorms:

- maximum t-conorm maxnorm (Def. 7),
- probabilistic sum probsum_conorm (Def. 9),
- bounded sum BoundedSum_conorm (Def. 19),
- drastic t-conorm drastic_conorm (Def. 14),
- nilpotent maximum nilmax_conorm (Def. 18),
- Hamacher t-conorm Hamacher_conorm (Def. 17).

Their basic properties and duality are shown; we also proved the predicate of the ordering of norms [10], [9]. It was proven formally that drastic-norm is the pointwise smallest t-norm and minnorm is the pointwise largest t-norm (maxnorm is the pointwise smallest t-conorm and drastic-conorm is the pointwise largest t-conorm). ADAM GRABOWSKI

This work is a continuation of the development of fuzzy sets in Mizar [6] started in [11] and [3]; it could be used to give a variety of more general operations on fuzzy sets. Our formalization is much closer to the set theory used within the Mizar Mathematical Library than the development of rough sets [4], the approach which was chosen allows however for merging both theories [5], [7].

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1. Preliminaries

One can verify that [0, 1] is non empty.

Let us consider elements a, b of [0, 1]. Now we state the propositions:

- (1) $\min(a, b) \in [0, 1].$
- (2) $\max(a, b) \in [0, 1].$
- $(3) \quad a \cdot b \in [0,1].$
- (4) $\max(0, a + b 1) \in [0, 1].$
- (5) $\min(a+b,1) \in [0,1].$
- (6) Let us consider elements a, b, c of [0, 1]. Then $\max(0, \max(0, a+b-1) + c-1) = \max(0, a + \max(0, b+c-1) 1)$.
- (7) Let us consider an element a of [0, 1]. Then $1 a \in [0, 1]$.

Let us consider elements a, b of [0, 1]. Now we state the propositions:

- (8) $a+b-(a \cdot b) \in [0,1]$. The theorem is a consequence of (7) and (3).
- (9) $\frac{a \cdot b}{a+b-(a \cdot b)} \in [0, 1]$. The theorem is a consequence of (3) and (8).
- (10) If $\max(a, b) \neq 1$, then $a \neq 1$ and $b \neq 1$.
- (11) Let us consider elements x, y of [0, 1]. If $x \cdot y = x + y$, then x = 0. The theorem is a consequence of (7).

Let us consider elements a, b of [0, 1]. Now we state the propositions:

- (12) $\max(a,b) = 1 \min(1-a,1-b).$
- (13) $\min(a+b,1) = 1 \max(0,1-a+(1-b)-1).$
- (14) $\frac{a+b-(2\cdot a\cdot b)}{1-(a\cdot b)} \in [0,1]$. The theorem is a consequence of (7) and (3).

Let f be a binary operation on [0, 1] and a, b be real numbers. Let us observe that f(a, b) is real.

Now we state the propositions:

(15) Let us consider real numbers a, b, and a binary operation t on [0, 1]. Then $t(a, b) \in [0, 1]$.

- (16) Let us consider a binary operation f on [0,1], and real numbers a, b. Then $1 - f(1 - a, 1 - b) \in [0, 1]$. The theorem is a consequence of (15) and (7).
- (17) Let us consider real numbers x, y, k. Suppose $k \leq 0$. Then
 - (i) $k \cdot \min(x, y) = \max(k \cdot x, k \cdot y)$, and
 - (ii) $k \cdot \max(x, y) = \min(k \cdot x, k \cdot y).$

2. BASIC EXAMPLE OF A TRIANGULAR NORM AND CONORM: MIN AND MAX

Let A be a real-membered set and f be a binary operation on A. We say that f is monotonic if and only if

(Def. 1) for every elements a, b, c, d of A such that $a \leq c$ and $b \leq d$ holds $f(a, b) \leq f(c, d)$.

We say that f has 1-identity if and only if

- (Def. 2) for every element a of A, f(a, 1) = a. We say that f has 1-annihilating if and only if
- (Def. 3) for every element a of A, f(a, 1) = 1. We say that f has 0-identity if and only if
- (Def. 4) for every element a of A, f(a, 0) = a.

We say that f has 0-annihilating if and only if

(Def. 5) for every element a of A, f(a, 0) = 0.

The scheme ExBinOp deals with a non empty, real-membered set \mathcal{A} and a binary functor \mathcal{F} yielding a set and states that

(Sch. 1) There exists a binary operation f on \mathcal{A} such that for every elements a, b of \mathcal{A} , $f(a,b) = \mathcal{F}(a,b)$

provided

• for every elements a, b of $\mathcal{A}, \mathcal{F}(a, b) \in \mathcal{A}$.

The functor minnorm yielding a binary operation on [0, 1] is defined by (Def. 6) for every elements a, b of $[0, 1], it(a, b) = \min(a, b)$.

Observe that minnorm is commutative, associative, and monotonic and has 1-identity and there exists a binary operation on [0, 1] which is commutative, associative, and monotonic and has 1-identity.

A t-norm is a commutative, associative, monotonic binary operation on [0, 1] with 1-identity. The functor maxnorm yielding a binary operation on [0, 1] is defined by

(Def. 7) for every elements a, b of $[0, 1], it(a, b) = \max(a, b)$.

One can verify that maxnorm is commutative, associative, and monotonic and has 0-identity and there exists a binary operation on [0, 1] which is commutative, associative, and monotonic and has 0-identity.

A t-conorm is a commutative, associative, monotonic binary operation on [0, 1] with 0-identity. Now we state the propositions:

- (18) Let us consider a commutative, monotonic binary operation t on [0, 1] with 1-identity, and an element a of [0, 1]. Then t(a, 0) = 0. The theorem is a consequence of (15).
- (19) Let us consider a commutative, monotonic binary operation t on [0, 1] with 0-identity, and an element a of [0, 1]. Then t(a, 1) = 1. The theorem is a consequence of (15).

Let us note that every commutative, monotonic binary operation on [0, 1] with 1-identity has 0-annihilating and every commutative, monotonic binary operation on [0, 1] with 0-identity has 1-annihilating.

3. Further Examples of Triangular Norms

The functor prodnorm yielding a binary operation on [0, 1] is defined by

(Def. 8) for every elements a, b of $[0, 1], it(a, b) = a \cdot b$.

Let us observe that prodnorm is commutative, associative, and monotonic and has 1-identity.

The functor problum-conorm yielding a binary operation on [0, 1] is defined by

(Def. 9) for every elements a, b of $[0, 1], it(a, b) = a + b - (a \cdot b)$.

The functor Lukasiewicz-norm yielding a binary operation on [0, 1] is defined by

(Def. 10) for every elements a, b of $[0, 1], it(a, b) = \max(0, a + b - 1)$.

One can check that Lukasiewicz-norm is commutative, associative, and monotonic and has 1-identity.

The functor drastic-norm yielding a binary operation on [0, 1] is defined by (Def. 11) for every elements a, b of [0, 1], if $\max(a, b) = 1$, then $it(a, b) = \min(a, b)$ and if $\max(a, b) \neq 1$, then it(a, b) = 0.

Now we state the proposition:

(20) Let us consider elements a, b of [0, 1]. Then

(i) if a = 1, then (drastic-norm)(a, b) = b, and

(ii) if b = 1, then (drastic-norm)(a, b) = a, and

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(iii) if $a \neq 1$ and $b \neq 1$, then (drastic-norm)(a, b) = 0.

Note that drastic-norm is commutative, associative, and monotonic and has 1-identity.

The functor nilmin-norm yielding a binary operation on [0, 1] is defined by

(Def. 12) for every elements a, b of [0, 1], if a + b > 1, then $it(a, b) = \min(a, b)$ and if $a + b \leq 1$, then it(a, b) = 0.

Observe that nilmin-norm is commutative, associative, and monotonic and has 1-identity.

The functor Hamacher-norm yielding a binary operation on [0, 1] is defined by

(Def. 13) for every elements a, b of $[0, 1], it(a, b) = \frac{a \cdot b}{a + b - (a \cdot b)}$.

One can verify that Hamacher-norm is commutative, associative, and monotonic and has 1-identity.

4. Basic Triangular Conorms

The functor drastic-conorm yielding a binary operation on [0, 1] is defined by

(Def. 14) for every elements a, b of [0, 1], if $\min(a, b) = 0$, then $it(a, b) = \max(a, b)$ and if $\min(a, b) \neq 0$, then it(a, b) = 1.

5. TRANSLATING BETWEEN TRIANGULAR NORMS AND CONORMS

Let t be a binary operation on [0, 1]. The functor conorm t yielding a binary operation on [0, 1] is defined by

(Def. 15) for every elements a, b of [0, 1], it(a, b) = 1 - t(1 - a, 1 - b).

Let t be a t-norm. Let us observe that conorm t is monotonic, commutative, and associative and has 0-identity.

Now we state the propositions:

- (21) maxnorm = conorm minnorm. PROOF: For every elements a, b of [0, 1], (maxnorm)(a, b) = 1 - (minnorm)(1 - a, 1 - b) by (7), (17), [12, (42)]. \Box
- (22) Let us consider a binary operation t on [0, 1]. Then conorm conorm t = t. The theorem is a consequence of (7).

6. The Ordering of Triangular Norms (and Conorms)

Let f_1, f_2 be binary operations on [0, 1]. We say that $f_1 \leq f_2$ if and only if (Def. 16) for every elements a, b of $[0, 1], f_1(a, b) \leq f_2(a, b)$.

Let us consider a t-norm t. Now we state the propositions:

- (23) drastic-norm $\leq t$. The theorem is a consequence of (20).
- (24) $t \leq \text{minnorm.}$

Now we state the proposition:

(25) Let us consider t-norms t_1, t_2 . If $t_1 \leq t_2$, then conorm $t_2 \leq \text{conorm} t_1$. The theorem is a consequence of (7).

7. TRIANGULAR CONORMS GENERATED FROM T-NORMS

The functor Hamacher-conorm yielding a binary operation on [0, 1] is defined by

(Def. 17) for every elements a, b of [0, 1], if a = b = 1, then it(a, b) = 1 and if $a \neq 1$ or $b \neq 1$, then $it(a, b) = \frac{a+b-(2\cdot a\cdot b)}{1-(a\cdot b)}$.

Now we state the proposition:

(26) conorm Hamacher-norm = Hamacher-conorm. The theorem is a consequence of (7).

Let us note that Hamacher-conorm is commutative, associative, and monotonic and has 0-identity.

Now we state the propositions:

- (27) conorm drastic-norm = drastic-conorm. The theorem is a consequence of (7).
- (28) conorm product problem (28) (7).

One can check that probsum-conorm is commutative, associative, and monotonic and has 0-identity.

The functor nilmax-conorm yielding a binary operation on [0, 1] is defined by

(Def. 18) for every elements a, b of [0, 1], if a + b < 1, then $it(a, b) = \max(a, b)$ and if $a + b \ge 1$, then it(a, b) = 1.

Now we state the proposition:

(29) conorm nilmin-norm = nilmax-conorm. The theorem is a consequence of (7) and (12).

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Let us note that nilmax-conorm is commutative, associative, and monotonic and has 0-identity.

The functor BoundedSum-conorm yielding a binary operation on [0, 1] is defined by

(Def. 19) for every elements a, b of $[0, 1], it(a, b) = \min(a + b, 1)$.

Now we state the proposition:

(30) conorm Lukasiewicz-norm = BoundedSum-conorm. The theorem is a consequence of (7) and (13).

One can check that BoundedSum-conorm is commutative, associative, and monotonic and has 0-identity.

Let us consider a t-conorm t. Now we state the propositions:

- (31) maxnorm $\leq t$.
- (32) $t \leq \text{drastic-conorm.}$

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