High-Speed Algorithms for RSA Cryptograms

Yasushi Fuwa Yoshinori Fujisawa Shinshu University Nagano Nagano

Summary. In this article, we propose a new high-speed processing method for encoding and decoding the RSA cryptogram that is a kind of public-key cryptogram. This cryptogram is not only used for encrypting data, but also for such purposes as authentication. However, the encoding and decoding processes take a long time because they require a great deal of calculations. As a result, this cryptogram is not suited for practical use. Until now, we proposed a highspeed algorithm of addition using $radix-2^k$ signed-digit numbers and clarified correctness of it ([5]). In this article, we defined two new operations for a highspeed coding and encoding processes on public-key cryptograms based on radix- 2^k signed-digit (SD) numbers. One is calculation of $(a*b) \mod c$ (a, b, c are natural)numbers). Another one is calculation of $(a^b) \mod c$ (a, b, c are natural numbers). Their calculations are realized repetition of addition. We propose a high-speed algorithm of their calculations using proposed addition algorithm and clarify the correctness of them. In the first section, we prepared some useful theorems for natural numbers and integers and so on. In the second section, we proved some properties of addition operation using a radix- 2^k SD numbers. In the third section, we defined some functions on the relation between a finite sequence of k-SD and a finite sequence of $\mathbb N$ and proved some properties about them. In the fourth section, algorithm of calculation of $(a * b) \mod c$ based on radix-2^k SD numbers is proposed and its correctness is clarified. In the last section, algorithm of calculation of $(a^b) \mod c$ based on radix- 2^k SD numbers is proposed and we clarified its correctness.

MML Identifier: RADIX_2.

The notation and terminology used in this paper are introduced in the following articles: [8], [6], [2], [3], [4], [9], [1], [5], [10], and [7].

C 2001 University of Białystok ISSN 1426-2630

1. Some Useful Theorems

In this paper k is a natural number.

The following propositions are true:

- (1) For every natural number a holds $a \mod 1 = 0$.
- (2) Let a, b be integers and n be a natural number. If n > 0, then $((a \mod n) + (b \mod n)) \mod n = (a + (b \mod n)) \mod n$ and $((a \mod n) + (b \mod n)) \mod n = ((a \mod n) + b) \mod n$.
- (3) For all integers a, b and for every natural number n such that n > 0holds $a \cdot b \mod n = a \cdot (b \mod n) \mod n$ and $a \cdot b \mod n = (a \mod n) \cdot b \mod n$.
- (4) For all natural numbers a, b, i such that $1 \leq i$ and 0 < b holds $(a \mod b_{\mathbb{N}}^{i}) \div b_{\mathbb{N}}^{i-1} = (a \div b_{\mathbb{N}}^{i-1}) \mod b.$
- (5) For all natural numbers i, n such that $i \in \text{Seg } n$ holds $i + 1 \in \text{Seg}(n+1)$.

2. Properties of Addition Operation Using Radix- 2^k Signed-Digit Numbers

One can prove the following propositions:

- (6) For every natural number k holds $\operatorname{Radix} k > 0$.
- (7) For every tuple x of 1 and k-SD holds SDDec x = DigA(x, 1).
- (8) For every integer x holds $SD_Add_Data(x, k) + SD_Add_Carry x \cdot Radix k = x.$
- (9) Let n be a natural number, x be a tuple of n + 1 and k SD, and x_1 be a tuple of n and k SD. Suppose that for every natural number i such that $i \in \text{Seg } n$ holds $x(i) = x_1(i)$. Then $\sum \text{DigitSD} x = \sum ((\text{DigitSD} x_1) \land (\text{SubDigit}(x, n + 1, k))).$
- (10) Let n be a natural number, x be a tuple of n+1 and k-SD, and x_1 be a tuple of n and k-SD. Suppose that for every natural number i such that $i \in \text{Seg } n$ holds $x(i) = x_1(i)$. Then $\text{SDDec } x_1 + ((\text{Radix } k)^n_{\mathbb{N}}) \cdot \text{DigA}(x, n+1) = \text{SDDec } x$.
- (11) Let n be a natural number. Suppose $n \ge 1$. Let x, y be tuples of n and k-SD. If $k \ge 2$, then SDDec x' + y +SD_Add_Carry DigA(x, n) +DigA $(y, n) \cdot ($ Radix $k)_{\mathbb{N}}^n =$ SDDec x +SDDec y.

3. Definitions on the Relation Between a Finite Sequence of k-SD and a Finite Sequence of \mathbb{N} and Some Properties about them

Let i, k, n be natural numbers and let x be a tuple of n and \mathbb{N} . The functor SubDigit2(x, i, k) yielding an element of \mathbb{N} is defined by:

(Def. 1) SubDigit2 $(x, i, k) = ((\operatorname{Radix} k)_{\mathbb{N}}^{i-1}) \cdot x(i).$

Let n, k be natural numbers and let x be a tuple of n and \mathbb{N} . The functor DigitSD2(x, k) yields a tuple of n and \mathbb{N} and is defined as follows:

(Def. 2) For every natural number i such that $i \in \text{Seg } n$ holds $\pi_i \text{DigitSD2}(x, k) = \text{SubDigit2}(x, i, k)$.

Let n, k be natural numbers and let x be a tuple of n and N. The functor SDDec2(x,k) yielding a natural number is defined as follows:

(Def. 3) SDDec2 $(x, k) = \sum DigitSD2(x, k)$.

Let i, k, x be natural numbers. The functor DigitDC2(x, i, k) yields a natural number and is defined as follows:

- (Def. 4) DigitDC2 $(x, i, k) = (x \mod (\operatorname{Radix} k)^i_{\mathbb{N}}) \div (\operatorname{Radix} k)^{i-1}_{\mathbb{N}}$. Let k, n, x be natural numbers. The functor DecSD2(x, n, k) yielding a tuple of n and \mathbb{N} is defined by:
- (Def. 5) For every natural number i such that $i \in \text{Seg } n$ holds (DecSD2(x, n, k))(i) = DigitDC2(x, i, k).

The following propositions are true:

- (12) Let n, k be natural numbers, x be a tuple of n and N, and y be a tuple of n and k-SD. If x = y, then DigitSD2(x, k) =DigitSD y.
- (13) Let n, k be natural numbers, x be a tuple of n and N, and y be a tuple of n and k-SD. If x = y, then SDDec2(x, k) = SDDec y.
- (14) For all natural numbers x, n, k holds DecSD2(x, n, k) = DecSD(x, n, k).
- (15) Let n be a natural number. Suppose $n \ge 1$. Let m, k be natural numbers. If m is represented by n, k, then m = SDDec2(DecSD2(m, n, k), k).
- 4. A High-Speed Algorithm of Calculation of $(a * b) \mod b$ Based on Radix-2^k Signed-Digit Numbers and its Correctness

Let q be an integer, let f, j, k, n be natural numbers, and let c be a tuple of n and k-SD. The functor Table1(q, c, f, j) yielding an integer is defined as follows:

(Def. 6) Table1 $(q, c, f, j) = q \cdot \text{DigA}(c, j) \mod f$.

Let q be an integer, let k, f, n be natural numbers, and let c be a tuple of n and k-SD. Let us assume that $n \ge 1$. The functor Mul_mod(q, c, f, k) yielding a tuple of n and \mathbb{Z} is defined by the conditions (Def. 7).

(Def. 7)(i) (Mul_mod(q, c, f, k))(1) = Table1(q, c, f, n), and

(ii) for every natural number i such that $1 \leq i$ and $i \leq n-1$ there exist integers I_1 , I_2 such that $I_1 = (\text{Mul}_\text{mod}(q, c, f, k))(i)$ and $I_2 = (\text{Mul}_\text{mod}(q, c, f, k))(i+1)$ and $I_2 = (\text{Radix } k \cdot I_1 + \text{Table1}(q, c, f, n - i)) \mod f$.

One can prove the following proposition

- (16) Let n be a natural number. Suppose $n \ge 1$. Let q be an integer and i_1, f, k be natural numbers. Suppose i_1 is represented by n, k and f > 0. Let c be a tuple of n and k-SD. If $c = \text{DecSD}(i_1, n, k)$, then $(\text{Mul}_mod(q, c, f, k))(n) = q \cdot i_1 \mod f$.
- 5. A High-Speed Algorithm of Calculation of (a^b) mod b Based on a Radix- 2^k Signed-Digit Numbers and its Correctness

Let n, f, j, m be natural numbers and let e be a tuple of n and \mathbb{N} . The functor Table2(m, e, f, j) yielding a natural number is defined as follows:

(Def. 8) Table2 $(m, e, f, j) = (m_{\mathbb{N}}^{\pi_j e}) \mod f$.

Let k, f, m, n be natural numbers and let e be a tuple of n and \mathbb{N} . Let us assume that $n \ge 1$. The functor Pow_mod(m, e, f, k) yields a tuple of n and \mathbb{N} and is defined by the conditions (Def. 9).

(Def. 9)(i) (Pow_mod(m, e, f, k))(1) = Table2(m, e, f, n), and

(ii) for every natural number i such that $1 \leq i$ and $i \leq n-1$ there exist natural numbers i_2 , i_3 such that $i_2 = (\text{Pow}_\text{mod}(m, e, f, k))(i)$ and $i_3 = (\text{Pow}_\text{mod}(m, e, f, k))(i+1)$ and $i_3 = ((i_2 \mathbb{N}^{\text{Radix } k}) \mod f) \cdot \text{Table2}(m, e, f, n-i) \mod f$.

One can prove the following proposition

(17) Let n be a natural number. Suppose $n \ge 1$. Let m, k, f, i_4 be natural numbers. Suppose i_4 is represented by n, k and f > 0. Let e be a tuple of n and \mathbb{N} . If $e = \text{DecSD2}(i_4, n, k)$, then $(\text{Pow}_mod(m, e, f, k))(n) = (m_{\mathbb{N}}^{i_4}) \mod f$.

References

- Grzegorz Bancerek. The fundamental properties of natural numbers. Formalized Mathematics, 1(1):41-46, 1990.
- Grzegorz Bancerek and Krzysztof Hryniewiecki. Segments of natural numbers and finite sequences. Formalized Mathematics, 1(1):107–114, 1990.

278

- [3] Czesław Byliński. Finite sequences and tuples of elements of a non-empty sets. *Formalized Mathematics*, 1(3):529–536, 1990.
- [4] Czesław Byliński. Functions and their basic properties. *Formalized Mathematics*, 1(1):55–65, 1990.
- [5] Yoshinori Fujisawa and Yasushi Fuwa. Definitions of radix- 2^k signed-digit number and its adder algorithm. *Formalized Mathematics*, 9(1):71–75, 2001.
- [6] Yoshinori Fujisawa, Yasushi Fuwa, and Hidetaka Shimizu. Euler's Theorem and small Fermat's Theorem. *Formalized Mathematics*, 7(1):123–126, 1998.
- [7] Andrzej Kondracki. The Chinese Remainder Theorem. Formalized Mathematics, 6(4):573–577, 1997.
- [8] Takaya Nishiyama and Yasuho Mizuhara. Binary arithmetics. Formalized Mathematics, 4(1):83–86, 1993.
- [9] Michał J. Trybulec. Integers. Formalized Mathematics, 1(3):501–505, 1990.
- [10] Zinaida Trybulec. Properties of subsets. Formalized Mathematics, 1(1):67–71, 1990.

Received February 1, 2000