

# Semantics of MML Query - Ordering

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**Summary.** Semantics of order directives of MML Query is presented. The formalization is done according to [1].

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The notation and terminology used in this paper have been introduced in the following articles: [2], [7], [13], [9], [10], [8], [3], [4], [5], [11], [17], [19], [18], [6], [15], [16], [14], and [12].

# 1. Preliminaries

In this paper X denotes a set, R,  $R_1$ ,  $R_2$  denote binary relations, x, y, z denote sets, and n, m, k denote natural numbers.

Let us consider a binary relation R on X. Now we state the propositions:

- (1) field  $R \subseteq X$ .
- (2) If  $x, y \in R$ , then  $x, y \in X$ .

Now we state the propositions:

- (3) Let us consider sets X, Y. Then  $(id_X)^{\circ}Y = X \cap Y$ .
- (4)  $\langle x, y \rangle \in \mathbb{R}$  |  $^2X$  if and only if  $x, y \in X$  and  $\langle x, y \rangle \in \mathbb{R}$ .
- (5)  $\operatorname{dom}(X \upharpoonright R) \subseteq \operatorname{dom} R$ .
- (6) Let us consider a total reflexive binary relation R on X and a subset S of X. Then  $R \mid^2 S$  is a total reflexive binary relation on S. The theorem is a consequence of (4). PROOF: Set  $Q = R \mid^2 S$ . dom Q = S.  $\square$
- (7) Let us consider transfinite sequences f, g. Then  $\operatorname{rng}(f \cap g) = \operatorname{rng} f \cup \operatorname{rng} g$ . Let us consider R. Let us note that R is transitive if and only if the condition (Def. 1) is satisfied.

(Def. 1) If  $x, y \in R$  and  $y, z \in R$ , then  $x, z \in R$ .

One can verify that R is antisymmetric if and only if the condition (Def. 2) is satisfied.

(Def. 2) If  $x, y \in R$  and  $y, x \in R$ , then x = y.

Now we state the proposition:

(8) Let us consider a non empty set X, a total connected binary relation R on X, and elements x, y of X. If  $x \neq y$ , then  $x, y \in R$  or  $y, x \in R$ .

# 2. Composition of Orders

Let  $R_1$ ,  $R_2$  be binary relations. The functor  $R_1$ ,  $R_2$  yielding a binary relation is defined by the term

(Def. 3)  $R_1 \cup (R_2 \setminus R_1^{\smile})$ .

Now we state the propositions:

- (9)  $x, y \in R_1, R_2$  if and only if  $x, y \in R_1$  or  $y, x \notin R_1$  and  $x, y \in R_2$ .
- (10) field  $(R_1, R_2)$  = field  $R_1 \cup$  field  $R_2$ . The theorem is a consequence of (9).
- (11)  $R_1, R_2 \subseteq R_1 \cup R_2$ . The theorem is a consequence of (9).

Let X be a set and  $R_1$ ,  $R_2$  be binary relations on X. Note that the functor  $R_1$ ,  $R_2$  yields a binary relation on X. Let  $R_1$ ,  $R_2$  be reflexive binary relations. One can verify that  $R_1$ ,  $R_2$  is reflexive.

Let  $R_1$ ,  $R_2$  be antisymmetric binary relations. Note that  $R_1$ ,  $R_2$  is antisymmetric.

Let X be a set and R be a binary relation on X. We say that R is  $\beta$ -transitive if and only if

(Def. 4) Let us consider elements x, y of X. If  $x, y \notin R$ , then for every element z of X such that  $x, z \in R$  holds  $y, z \in R$ .

Observe that every binary relation on X which is connected total and transitive is also  $\beta$ -transitive.

Let us observe that there exists an order in X which is connected.

Let  $R_1$  be a  $\beta$ -transitive transitive binary relation on X and  $R_2$  be a transitive binary relation on X. Observe that  $R_1$ ,  $R_2$  is transitive.

Let  $R_1$  be a binary relation on X and  $R_2$  be a total reflexive binary relation on X. Let us note that  $R_1$ ,  $R_2$  is total and reflexive as a binary relation on X.

Let  $R_2$  be a total connected reflexive binary relation on X. One can verify that  $R_1$ ,  $R_2$  is connected.

Now we state the propositions:

- (12)  $(R, R_1), R_2 = R, (R_1, R_2)$ . The theorem is a consequence of (9).
- (13) Let us consider a connected reflexive total binary relation R on X and a binary relation  $R_2$  on X. Then R,  $R_2 = R$ . The theorem is a consequence of (9) and (2).

#### 3. number of ORDERING

Let X be a set and f be a function from X into  $\mathbb{N}$ . The functor number of f yielding a binary relation on X is defined by

(Def. 5)  $x, y \in it$  if and only if  $x, y \in X$  and f(x) < f(y).

Let us note that number of f is antisymmetric transitive and  $\beta$ -transitive.

Let X be a finite set and O be an operation of X. The functor value of O yielding a function from X into  $\mathbb{N}$  is defined by

(Def. 6) Let us consider an element x of X. Then  $it(x) = \overline{\overline{x(O)}}$ .

Now we state the proposition:

(14) Let us consider a finite set X, an operation O of X, and elements x, y of X. Then  $x, y \in$  number of value of O if and only if  $\overline{x(O)} < \overline{y(O)}$ .

Let us consider X. Let O be an operation of X. The functor first O yielding a binary relation on X is defined by

(Def. 7) Let us consider elements x, y of X. Then  $x, y \in it$  if and only if  $x(O) \neq \emptyset$  and  $y(O) = \emptyset$ .

Let us observe that first O is antisymmetric transitive and  $\beta$ -transitive.

## 4. Ordering by Resources

Let A be a finite sequence and x be an element. The functor  $A \leftarrow x$  yielding a set is defined by the term

(Def. 8)  $\bigcap (A^{-1}(\{x\})).$ 

Let us consider x. Note that  $A \leftarrow x$  is natural.

Let us consider a finite sequence A. Now we state the propositions:

- (15) If  $x \notin \operatorname{rng} A$ , then  $A \leftarrow x = 0$ .
- (16) If  $x \in \operatorname{rng} A$ , then  $A \leftarrow x \in \operatorname{dom} A$  and  $x = A(A \leftarrow x)$ .
- (17) If  $A \leftarrow x = 0$ , then  $x \notin \operatorname{rng} A$ .

Let us consider X. Let A be a finite sequence and f be a function. The functor resource(X, A, f) yielding a binary relation on X is defined by

(Def. 9)  $x, y \in it$  if and only if  $x, y \in X$  and  $A \leftarrow (f(x)) \neq 0$  and  $A \leftarrow (f(x)) < A \leftarrow (f(y))$  or  $A \leftarrow (f(y)) = 0$ .

Let us observe that resource(X, A, f) is antisymmetric transitive and  $\beta$ -transitive.

#### 5. Ordering by Number of Iteration

Let us consider X. Let R be a binary relation on X and n be a natural number. One can check that the functor  $R^n$  yields a binary relation on X. Now we state the propositions:

- (18) If  $(R^n)^{\circ}X = \emptyset$  and  $m \ge n$ , then  $(R^m)^{\circ}X = \emptyset$ .
- (19) If for every n,  $(R^n)^{\circ}X \neq \emptyset$  and X is finite, then there exists x such that  $x \in X$  and for every n,  $(R^n)^{\circ}x \neq \emptyset$ . The theorem is a consequence of (18). PROOF: Define  $\mathcal{P}[\text{element}, \text{element}] \equiv \text{there exists } n \text{ such that } \$_2 = n \text{ and } (R^n)^{\circ}\$_1 = \emptyset$ . For every element x such that  $x \in X$  there exists an element y such that  $y \in \mathbb{N}$  and  $\mathcal{P}[x,y]$ . Consider f being a function such that  $\text{dom } f = X \text{ and rng } f \subseteq \mathbb{N} \text{ and for every element } x \text{ such that } x \in X \text{ holds } \mathcal{P}[x,f(x)]$ . Consider n such that  $\text{rng } f \subseteq \mathbb{Z}_n$ .  $\{\{x\} \text{ where } x \text{ is an element of } X: x \in X\} \subseteq 2^X$ . Reconsider  $Y = \{\{x\} \text{ where } x \text{ is an element of } X: x \in X\}$  as a family of subsets of X.  $X = \bigcup Y$ .  $\{(R^n)^{\circ}y \text{ where } y \text{ is a subset of } X: y \in Y\} \subseteq \{\emptyset\}$ .  $\square$
- (20) If R is reversely well founded and irreflexive and X is finite and R is finite, then there exists n such that  $(R^n)^{\circ}X = \emptyset$ . The theorem is a consequence of (19). PROOF: Define  $\mathcal{Q}[\text{element}] \equiv \text{for every } n, (R^n)^{\circ}\$_1 \neq \emptyset$ . Consider x0 being a set such that  $x0 \in X$  and  $\mathcal{Q}[x0]$ . Define  $\mathcal{P}[\text{element, element, element}] \equiv \text{if } \mathcal{Q}[\$_2]$ , then  $\$_3 \in R^{\circ}\$_2$  and  $\mathcal{Q}[\$_3]$ . For every natural number n and for every set x, there exists a set y such that  $\mathcal{P}[n, x, y]$ . Consider f being a function such that dom  $f = \mathbb{N}$  and f(0) = x0 and for every natural number n,  $\mathcal{P}[n, f(n), f(n+1)]$ . Define  $\mathcal{R}[\text{natural number}] \equiv \mathcal{Q}[f(\$_1)]$ . rng  $f \subseteq \text{field } R$ . Consider z being an element such that  $z \in \text{rng } f$  and for every element x such that  $x \in \text{rng } f$  and  $z \neq x$  holds  $\langle z, x \rangle \notin R$ . Consider z being an element such that  $z \in \text{rng } f$  and  $z \in \mathcal{P}[x]$ .

Let us consider X. Let O be an operation of X. Assume O is reversely well founded, irreflexive, and finite. The functor iteration of O yielding a binary relation on X is defined by

- (Def. 10) There exists a function f from X into  $\mathbb{N}$  such that
  - (i) it = number of f, and
  - (ii) for every element x of X such that  $x \in X$  there exists n such that f(x) = n and  $x(O^n) \neq \emptyset$  or n = 0 and  $x(O^n) = \emptyset$  and  $x(O^{n+1}) = \emptyset$ .

Let us note that every binary relation which is empty is also irreflexive and reversely well founded.

Let us consider X. Let us note that there exists an operation of X which is empty.

Let O be a reversely well founded irreflexive finite operation of X. One can check that **iteration** of O is antisymmetric transitive and  $\beta$ -transitive.

#### 6. value of Ordering

Let X be a finite set. Let us observe that every order in X is well founded. Note that every connected order in X is well-ordering.

Let us consider X. Let R be a connected order in X and S be a finite subset of X. The functor  $\operatorname{order}(S, R)$  yielding a finite 0-sequence of X is defined by

- (Def. 11) (i) rng it = S, and
  - (ii) it is one-to-one, and
  - (iii) for every natural numbers i, j such that  $i, j \in \text{dom } it \text{ holds } i \leq j \text{ iff } it(i), it(j) \in R.$

Now we state the proposition:

(21) Let us consider finite subsets  $S_1$ ,  $S_2$  of X and a connected order R in X. Then  $\operatorname{order}(S_1 \cup S_2, R) = \operatorname{order}(S_1, R) \cap \operatorname{order}(S_2, R)$  if and only if for every x and y such that  $x \in S_1$  and  $y \in S_2$  holds  $x \neq y$  and  $x, y \in R$ . The theorem is a consequence of (7). PROOF: Set  $o1 = \operatorname{order}(S_1, R)$ . Set  $o2 = \operatorname{order}(S_2, R)$ .  $\operatorname{order}(S_1, R) \cap \operatorname{order}(S_2, R)$  is one-to-one.  $\square$ 

Let X be a finite set, O be an operation of X, and R be a connected order in X. The functor value of(O, R) yielding a binary relation on X is defined by

(Def. 12) Let us consider elements x, y of X. Then  $x, y \in it$  if and only if  $x(O) \neq \emptyset$  and  $y(O) = \emptyset$  or  $y(O) \neq \emptyset$  and  $(\operatorname{order}(x(O), R))_0, (\operatorname{order}(y(O), R))_0 \in R$  and  $(\operatorname{order}(x(O), R))_0 \neq (\operatorname{order}(y(O), R))_0$ .

Let  $R_1$  be a connected order in X. One can check that value of  $(O, R_1)$  is antisymmetric transitive and  $\beta$ -transitive.

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