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Introduction to Galois Theory

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Summary. This article is the second in a series of five articles formalizing the Fundamental Theorem of Galois Theory [10], [7], [8] using the Mizar formalism [1], [2], [6].

We start the actual formalization defining groups of automorphisms and fixed fields and proving some of their basic properties [10]. We also define conjugates for a group of automorphisms and prove that for algebraic elements $a \in E$, there is a bijection between $\operatorname{Aut}(F(a),F)$ and the roots of a's minimal polynomial in F(a) [5]. Finally we define Galois extensions as extensions E over F with $\operatorname{Fix}(E,\operatorname{Aut}(E,F))=F$ and show that the complex numbers are a Galois extension of the real numbers. We also consider finite fields and prove that a field E of order p^n is a Galois extension of \mathbb{Z}_p of degree n and that $\operatorname{Aut}(E,\mathbb{Z}_p)$ is generated by the Frobenius morphism [8].

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1. Preliminaries

Let X be a non empty set and x be an element of X. Observe that the functor $\{x\}$ yields a subset of X. Let F be a non quadratic complete field and a be a non square element of F. One can check that X^2 - a is irreducible and there exists an extension of F which is F-quadratic.

Let F be a field. Note that every extension of F which is F-quadratic is also F-simple.

Let E be an extension of F and a be an F-algebraic element of E. One can verify that $\text{Roots}(\text{FAdj}(F,\{a\}), \text{MinPoly}(a,F))$ is non empty and finite.

Now we state the proposition:

(1) Let us consider an element z of \mathbb{C}_{F} . Then z is an element of \mathbb{R}_{F} if and only if $\overline{z} = z$.

One can check that \mathbb{C}_F and has not characteristic 2.

One can check that the functor i yields a non zero, (\mathbb{R}_{F})-algebraic element of \mathbb{C}_{F} . Let us note that $X^{2}+1_{\mathbb{R}_{F}}$ is irreducible.

Now we state the propositions:

- (2) Roots($\mathbb{C}_{F}, X^{2} + 1_{\mathbb{R}_{F}}$) = $\{i, -i\}$.
- (3) MinPoly $(i, \mathbb{R}_F) = X^2 + 1_{\mathbb{R}_F}$. The theorem is a consequence of (2).
- (4) \mathbb{C}_F is a splitting field of $X^2 + 1_{\mathbb{R}_F}$. The theorem is a consequence of (2).
- (5) $\mathbb{C}_{F} = \text{FAdj}(\mathbb{R}_{F}, \{i\})$. The theorem is a consequence of (4) and (2).

One can verify that \mathbb{C}_F is (\mathbb{R}_F) -quadratic.

2. Groups of Automorphims

Let F be a field. The functor Auts(F) yielding a non empty set is defined by the term

(Def. 1) the set of all f where f is an automorphism of F.

One can check that Auts(F) is F-functional and every element of Auts(F) is additive, multiplicative, unity-preserving, and isomorphism.

The functor AutComp(F) yielding a binary operation on Auts(F) is defined by

(Def. 2) for every automorphisms f, g of F, $it(f, g) = f \cdot g$.

The functor $\mathrm{Aut}(F)$ yielding a strict, non empty multiplicative magma is defined by the term

(Def. 3) $\langle Auts(F), AutComp(F) \rangle$.

Let us note that Aut(F) is group-like and associative.

Now we state the proposition:

(6) Let us consider a field F. Then $\mathbf{1}_{\operatorname{Aut}(F)} = \operatorname{id}_F$.

Let F be a field. Note that Aut(F) is F-functional and the carrier of Aut(F) is F-functional and every subset of Aut(F) is F-functional.

Let G be a subgroup of $\operatorname{Aut}(F)$. Observe that the carrier of G is F-functional and every element of the carrier of $\operatorname{Aut}(F)$ is additive, multiplicative, unity-preserving, and isomorphism.

Let G be a non empty subset of Aut(F). Let us note that every element of G is additive, multiplicative, unity-preserving, and isomorphism.

Let G be a non empty subgroup of Aut(F). Note that every element of the carrier of G is additive, multiplicative, unity-preserving, and isomorphism.

Let E be an extension of F. The functor Auts(E, F) yielding a non empty subset of Auts(E) is defined by the term

(Def. 4) the set of all f where f is an F-fixing automorphism of E.

One can verify that Auts(E, F) is E-functional and every element of Auts(E, F) is F-fixing, additive, multiplicative, unity-preserving, and isomorphism.

The functor AutComp(E, F) yielding a binary operation on Auts(E, F) is defined by the term

(Def. 5) $\operatorname{AutComp}(E) \upharpoonright \operatorname{Auts}(E, F)$.

The functor Aut(E, F) yielding a strict multiplicative magma is defined by the term

(Def. 6) $\langle \text{Auts}(E, F), \text{AutComp}(E, F) \rangle$.

Let us note that $\operatorname{Aut}(E,F)$ is non empty and $\operatorname{Aut}(E,F)$ is group-like and associative.

One can verify that the functor $\operatorname{Aut}(E,F)$ yields a strict subgroup of $\operatorname{Aut}(E)$. Let us observe that $\operatorname{Aut}(E,F)$ is E-functional and the carrier of $\operatorname{Aut}(E,F)$ is E-functional and every subset of $\operatorname{Aut}(E,F)$ is E-functional and every element of the carrier of $\operatorname{Aut}(E,F)$ is F-fixing, additive, multiplicative, unity-preserving, and isomorphism.

Let G be a non empty subset of Aut(E, F). Let us observe that every element of G is F-fixing, additive, multiplicative, unity-preserving, and isomorphism.

Let G be a subgroup of $\operatorname{Aut}(E,F)$. One can check that every element of the carrier of G is F-fixing, additive, multiplicative, unity-preserving, and isomorphism.

Let E be a field and F be a subfield of E. The functor Aut(E,F) yielding a strict subgroup of Aut(E) is defined by the term

(Def. 7) Aut(FieldExt(E, F), F).

Let F be a field, E be an extension of F, and K be an intermediate field of E, F. The functor Aut(E,K) yielding a strict subgroup of Aut(E) is defined by the term

(Def. 8) Aut(FieldExt(E, K), K).

Now we state the proposition:

(7) Let us consider a field F, an extension E of F, an intermediate field K_1 of E, F, and a subfield K_2 of E. If $K_2 = K_1$, then $Aut(E, K_2) = Aut(E, K_1)$.

3. Conjugates

Let F be a field, G be a subgroup of Aut(F), and a be an element of F. The

functor Conjugates(a, G) yielding a non empty subset of F is defined by the term

(Def. 9) the set of all f(a) where f is an element of the carrier of G.

Let E be an extension of F and a be an element of E. The functor Conjugates(a) yielding a non empty subset of E is defined by the term

(Def. 10) Conjugates (a, Aut(E, F)).

Let G be a subgroup of Aut(F) and a be an element of F. We introduce the notation Conj(a, G) as a synonym of Conjugates(a, G).

Let E be an extension of F and a be an element of E. We introduce the notation Conj(a) as a synonym of Conjugates(a).

Let us consider a field F, an extension E of F, an F-algebraic element a of E, and an element b of E. Now we state the propositions:

- (8) Suppose $E \approx \operatorname{FAdj}(F, \{a\})$. Then if $b \in \operatorname{Roots}(\operatorname{FAdj}(F, \{a\}), \operatorname{MinPoly}(a, F))$, then $\operatorname{FAdj}(F, \{b\}) = \operatorname{FAdj}(F, \{a\})$.
- (9) Suppose $b \in \text{Roots}(\text{FAdj}(F, \{a\}), \text{MinPoly}(a, F))$. Then there exists an F-fixing automorphism f of $\text{FAdj}(F, \{a\})$ such that f(a) = b. The theorem is a consequence of (8).

Now we state the proposition:

(10) Let us consider a field F, an extension E of F, a subgroup G of Aut(E, F), and an F-algebraic element a of E. Then $Conj(a, G) \subseteq Roots(E, MinPoly(a, F))$.

Let F be a field, E be an extension of F, G be a subgroup of Aut(E, F), and e be an F-algebraic element of E. Note that Conj(e, G) is finite and Conj(e) is finite.

Let us consider a field F, an extension E of F, and an F-algebraic element a of E. Now we state the propositions:

- (11) If $E \approx \text{FAdj}(F, \{a\})$, then Conj(a) = Roots(E, MinPoly(a, F)). The theorem is a consequence of (10) and (9).
- (12) There exists a function f from Auts(FAdj(F, {a}), F) into Roots(FAdj(F, {a}), Min such that f is bijective. PROOF: Set $G = \text{Auts}(\text{FAdj}(F, \{a\}), F)$. Set $R = \text{Roots}(\text{FAdj}(F, \{a\}), \text{MinPoly}(a, F))$

PROOF: Set $G = \operatorname{Auts}(\operatorname{FAdj}(F, \{a\}), F)$. Set $R = \operatorname{Roots}(\operatorname{FAdj}(F, \{a\}), \operatorname{Min.} Define \mathcal{P}[\operatorname{object}, \operatorname{object}] \equiv \operatorname{there}$ exists an element f of G such that $\$_1 = f$ and $\$_2 = f(a)$. Consider h being a function from G into R such that for every object x such that $x \in G$ holds $\mathcal{P}[x, h(x)]$ from $[4, \operatorname{Sch.} 1]$. \square

Let F be a field, E be an extension of F, and a be an F-algebraic element of E. Observe that $\operatorname{Aut}(\operatorname{FAdj}(F,\{a\}),F)$ is finite.

Let S be a subset of E. The functor Auts(S) yielding a non empty subset of Auts(E, F) is defined by the term

(Def. 11) $\{f, \text{ where } f \text{ is an } F\text{-fixing automorphism of } E: f^{\circ}S = S\}.$

Let a be an element of E. The functor Auts(a) yielding a non empty subset of Auts(E, F) is defined by the term

(Def. 12) Auts($\{a\}$).

Now we state the proposition:

(13) Let us consider a field F, an extension E of F, and an element a of E. Then $\operatorname{Auts}(a) = \{f, \text{ where } f \text{ is an } F\text{-fixing automorphism of } E: f(a) = a\}.$

Let F be a field, E be an extension of F, and S be a subset of E. The functor AutComp(S) yielding a binary operation on Auts(S) is defined by the term

(Def. 13) $\operatorname{AutComp}(E, F) \upharpoonright \operatorname{Auts}(S)$.

The functor $\mathrm{Aut}(S)$ yielding a strict, non empty multiplicative magma is defined by the term

(Def. 14) $\langle \text{Auts}(S), \text{AutComp}(S) \rangle$.

Let S be a non empty subset of E. One can check that Aut(S) is group-like and associative.

Let us observe that the functor Aut(S) yields a strict subgroup of Aut(E, F). Let a be an element of E. The functor Aut(a) yielding a strict subgroup of Aut(E, F) is defined by the term

(Def. 15) $\operatorname{Aut}(\{a\})$.

Now we state the proposition:

(14) Let us consider a field F, an extension E of F, and an element a of E. Then $\overline{\overline{\text{Conj}(a)}} = |\bullet: \text{Aut}(a)|$.

PROOF: Set $H = \operatorname{Aut}(a)$. Set $G = \operatorname{Aut}(E, F)$. Auts $(a) = \{f, \text{ where } f \text{ is an } F\text{-fixing automorphism of } E: f(a) = a\}$ and the carrier of $H = \operatorname{Auts}(a)$. Define $\mathcal{P}[\operatorname{object}, \operatorname{object}] \equiv \text{ for every element } f$ of the carrier of G such that $\$_1 = f(a)$ holds $\$_2 = f \cdot H$. Consider h being a function from $\operatorname{Conj}(a)$ into the left cosets of H such that for every object x such that $x \in \operatorname{Conj}(a)$ holds $\mathcal{P}[x, h(x)]$ from $[4, \operatorname{Sch}. 1]$. rng $h = \operatorname{the left}$ cosets of H. \square

4. Fixed Fields

Let F be a field and G be a subgroup of $\operatorname{Aut}(F)$. The functor $\operatorname{Fixed-Elements}(F,G)$ yielding a non empty subset of F is defined by the term

(Def. 16) $\{a, \text{ where } a \text{ is an element of } F : \text{ for every function } f \text{ from } F \text{ into } F \text{ such that } f \in G \text{ holds } f(a) = a\}.$

We introduce the notation Fixed-El(F, G) as a synonym of Fixed-Elements(F, G).

One can check that Fixed-El(F, G) is inducing subfield.

The functor Fix(F,G) yielding a strict double loop structure is defined by (Def. 17) the carrier of it = Fixed-El(F,G) and the addition of $it = (the addition of F) \upharpoonright Fixed-El(F,G)$ and the multiplication of $it = (the multiplication of F) \upharpoonright Fixed-El(F,G)$ and $1_{it} = 1_F$ and $0_{it} = 0_F$.

One can check that Fix(F, G) is non degenerated and Fix(F, G) is Abelian, add-associative, right zeroed, and right complementable and Fix(F, G) is commutative, associative, well unital, distributive, and almost left invertible.

Let us note that the functor Fix(F,G) yields a strict subfield of F. Let E be an extension of F and G be a subgroup of Aut(E,F). Let us note that the functor Fix(E,G) yields an intermediate field of E, F.

5. Some Basic Properties

Let us consider a field F and an extension E of F. Now we state the propositions:

- (15) F is a subfield of Fix(E, Aut(E, F)).
- (16) Every intermediate field of E, F is a subfield of Fix(E, Aut(E, K)). Now we state the propositions:
- (17) Let us consider a field F. Then every subgroup of Aut(F) is a subgroup of Aut(F, Fix(F, G)).
- (18) Let us consider a field F, an extension E of F, and an intermediate field K of E, F. Then Aut(E, K) is a subgroup of Aut(E, F).
- (19) Let us consider a field F, an extension E of F, and intermediate fields K_1 , K_2 of E, F. Suppose K_1 is a subfield of K_2 . Then $Aut(E, K_2)$ is a subgroup of $Aut(E, K_1)$.
- (20) Let us consider a field F, and subgroups G_1 , G_2 of Aut(F). Suppose G_1 is a subgroup of G_2 . Then $Fix(F, G_2)$ is a subfield of $Fix(F, G_1)$.
- (21) Let us consider a field F, and an extension E of F. Then Aut(E, F) = Aut(E, Fix(E, Aut(E, F))). The theorem is a consequence of (17) and (15).
- (22) Let us consider a field F, and a subgroup G of Aut(F). Then Fix(F, G) = Fix(F, Aut(F, Fix(F, G))). The theorem is a consequence of (17), (20), and (15).

6. Galois Extensions

Let F be a field and E be an extension of F. We say that E is F-Galois if and only if

(Def. 18) $\operatorname{Fix}(E, \operatorname{Aut}(E, F)) \approx F$.

One can check that there exists an extension of F which is F-Galois.

A Galois extension of F is a F-Galois extension of F. One can verify that there exists a Galois extension of F which is F-finite.

Now we state the propositions:

- (23) Let us consider a field F, and an extension E of F. Then E is F-Galois if and only if there exists a subgroup G of $\operatorname{Aut}(E)$ such that $\operatorname{Fix}(E,G)\approx F$. The theorem is a consequence of (20) and (15).
- (24) Every field is a Galois extension of F.

The functor *' yielding a function from \mathbb{C}_F into \mathbb{C}_F is defined by

(Def. 19) for every element z of \mathbb{C}_{F} , $it(z) = \overline{z}$.

Let us note that *' is (\mathbb{R}_F) -fixing and isomorphism.

Let us note that the functor *' yields an element of the carrier of $Aut(\mathbb{C}_F, \mathbb{R}_F)$. Now we state the propositions:

- (25) $\operatorname{Auts}(\mathbb{C}_F, \mathbb{R}_F) = \{ \operatorname{id}_{\mathbb{C}_F}, *' \}.$ The theorem is a consequence of (5), (2), and (3).
- (26) $\operatorname{Aut}(\mathbb{C}_F, \mathbb{R}_F) = \operatorname{gr}(\{*'\})$. The theorem is a consequence of (25).

One can verify that $\operatorname{Aut}(\mathbb{C}_F, \mathbb{R}_F)$ is finite and cyclic.

Now we state the propositions:

- (27) order $\operatorname{Aut}(\mathbb{C}_F, \mathbb{R}_F) = 2$. The theorem is a consequence of (25).
- (28) $\operatorname{Aut}(\mathbb{C}_F, \mathbb{R}_F)$ and \mathbb{Z}_2^+ are isomorphic.

PROOF: Set $E = \mathbb{C}_{F}$. Set $F = \mathbb{R}_{F}$. Define $\mathcal{P}[\text{object}, \text{object}] \equiv \$_{1} = 0$ and $\$_{2} = \text{id}_{E}$ or $\$_{1} = 1$ and $\$_{2} = *'$. Consider f being a function from the carrier of \mathbb{Z}_{2}^{+} into the carrier of Aut(E, F) such that for every object x such that $x \in \text{the carrier of } \mathbb{Z}_{2}^{+}$ holds $\mathcal{P}[x, f(x)]$ from [4, Sch. 1]. $\text{id}_{\mathbb{C}_{F}} \neq *'$ by [9, (2)], [3, (31), (17)]. $\overline{\alpha} = \overline{\beta}$, where α is the carrier of \mathbb{Z}_{2}^{+} and β is the carrier of Aut(E, F). \square

Observe that \mathbb{C}_F is (\mathbb{R}_F) -Galois.

Now we state the proposition:

(29) \mathbb{C}_{F} is a Galois extension of \mathbb{R}_{F} .

Let p be a prime number, n be a non zero natural number, and F be a Galois field of p^n . Observe that the functor $\operatorname{Frob}(F)$ yields an element of the carrier of $\operatorname{Aut}(F,\mathbb{Z}/p)$. Let m be a natural number. One can verify that $(\operatorname{Frob}(F))^m$ is (\mathbb{Z}/p) -fixing and isomorphism.

Let us consider a prime number p, a non zero natural number n, and a Galois field F of p^n . Now we state the propositions:

(30) $\operatorname{Aut}(F, \mathbb{Z}/p) = \operatorname{Aut}(F).$

- (31) Auts $(F, \mathbb{Z}/p) = \{(\operatorname{Frob}(F))^m, \text{ where } m \text{ is a natural number } : 0 \leq m \leq n-1\}$. The theorem is a consequence of (30).
- (32) $\operatorname{Aut}(F, \mathbb{Z}/p) = \operatorname{gr}(\{\operatorname{Frob}(F)\}).$

Let p be a prime number, n be a non zero natural number, and F be a Galois field of p^n . Observe that $\operatorname{Aut}(F,\mathbb{Z}/p)$ is finite and cyclic.

Let us consider a prime number p, a non zero natural number n, and a Galois field F of p^n . Now we state the propositions:

- (33) order $\operatorname{Aut}(F, \mathbb{Z}/p) = n$.
- (34) $\operatorname{Aut}(F, \mathbb{Z}/p)$ and \mathbb{Z}_n^+ are isomorphic.

PROOF: Set $a = \operatorname{Frob}(F)$. Define $\mathcal{P}[\operatorname{object}, \operatorname{object}] \equiv \operatorname{there}$ exists a natural number m such that $\$_1 = m$ and $\$_2 = a^m$. Consider f being a function from the carrier of \mathbb{Z}_n^+ into the carrier of $\operatorname{Aut}(F, \mathbb{Z}/p)$ such that for every object x such that $x \in \operatorname{the}$ carrier of \mathbb{Z}_n^+ holds $\mathcal{P}[x, f(x)]$ from [4, Sch. 1]. \square

Let p be a prime number and n be a non zero natural number. Let us note that every Galois field of p^n is (\mathbb{Z}/p) -Galois.

Now we state the proposition:

(35) Let us consider a prime number p, and a non zero natural number n. Then every Galois field of p^n is a Galois extension of \mathbb{Z}/p .

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